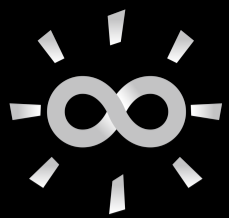


The Bright Side of Mathematics

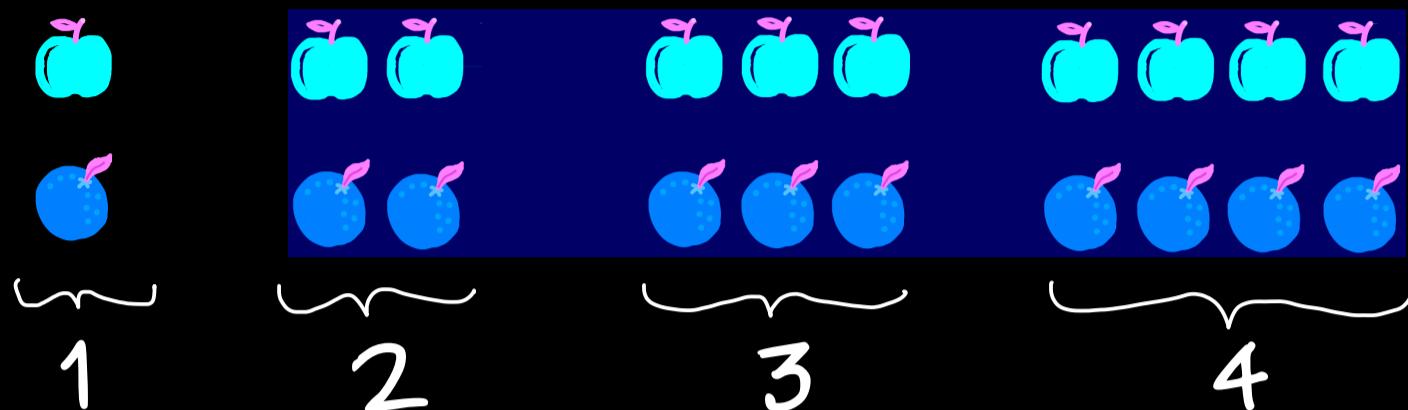
The following pages cover the whole Start Learning Numbers course of the Bright Side of Mathematics. Please note that the creator lives from generous supporters and would be very happy about a donation. See more here: <https://tbsom.de/support>

Have fun learning mathematics!



Start Learning Numbers - Part 1

Natural numbers



$$\mathbb{N} = \{1, 2, 3, 4, \dots\}$$

$$\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$$

$0 := \emptyset$ empty set

$1 := \{0\}$ set with one element

$2 := \{0, 1\}$ set with two elements

$3 := \{0, 1, 2\}$ set with three elements

$4 := \{0, 1, 2, 3\} = 3 \cup \{3\}$

\vdots

Axiom: There is a set \mathbb{N}_0 with the properties:

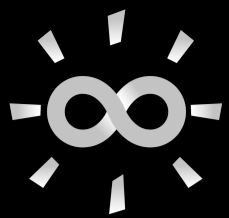
(a) $0 \in \mathbb{N}_0$

(b) $\forall x: x \in \mathbb{N}_0 \rightarrow x \cup \{x\} \in \mathbb{N}_0$

And \mathbb{N}_0 is the smallest set having these two properties.

Successor map:

$$s: \mathbb{N}_0 \rightarrow \mathbb{N}_0$$
$$x \mapsto x \cup \{x\}, \quad s(6) = 7$$



Start Learning Numbers - Part 2

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

Properties of \mathbb{N}_0 :

(1) $0 \in \mathbb{N}_0$

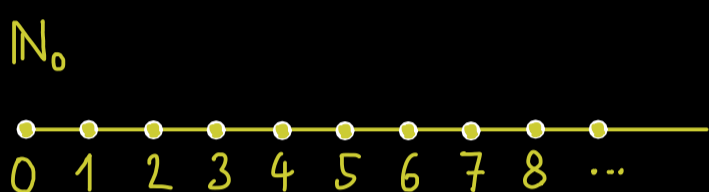
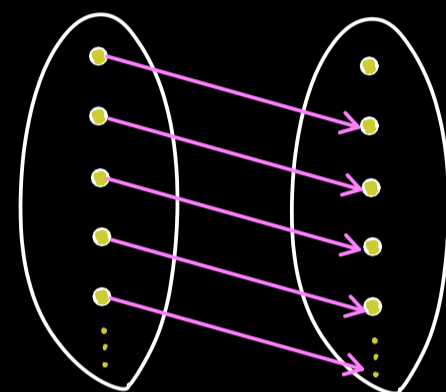
(2) There is a map $s: \mathbb{N}_0 \rightarrow \mathbb{N}_0$ that satisfies:

(2a) s is injective

(2b) $0 \notin \text{Ran}(s) = s[\mathbb{N}_0]$

(2c) If $M \subseteq \mathbb{N}_0$ with
 $0 \in M$ and $s[M] \subseteq M$,
then $M = \mathbb{N}_0$.

(mathematical induction)



Addition in \mathbb{N}_0 : map $\mathbb{N}_0 \times \mathbb{N}_0 \rightarrow \mathbb{N}_0$

$(m, n) \mapsto m + n$

How is it defined? $2 + 4 := 6$

$m + 0 := m$, $m + 1 := s(m)$, $m + 2 := s(m + 1)$

Recursive definition:

$m + s(n) := s(m + n)$

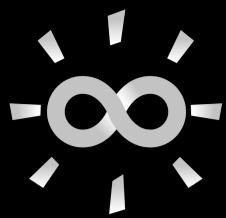
$2 + 5 = 2 + s(4) = s(2 + 4) = s(6) = 7$

Dedekind's principle of recursive definition:

For a set A , $a \in A$ and $h: A \rightarrow A$, then there exists a unique map

$f: \mathbb{N}_0 \rightarrow A$ with $f(0) = a$ and $f(s(n)) = h(f(n))$.

(" $a, h(a), h(h(a)), h(h(h(a))), \dots$ ")



Start Learning Numbers - Part 3

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

Each $n \in \mathbb{N}_0$ has a unique successor:

$$s: \mathbb{N}_0 \rightarrow \mathbb{N}_0, \quad s(n) = n + 1$$

We already know: $m + (n + 1) = (m + n) + 1$ (RD)

Mathematical induction:

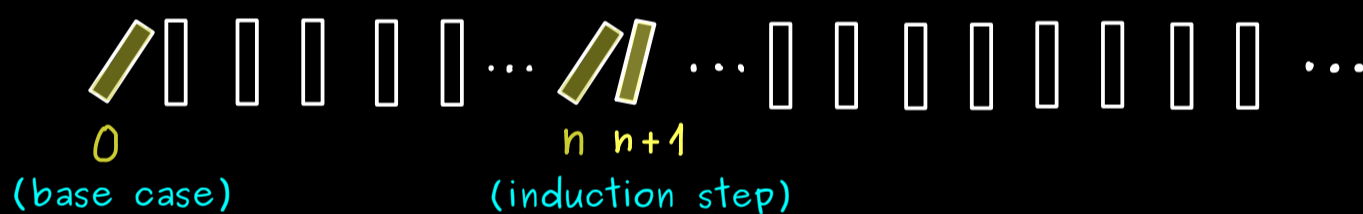
\mathbb{N}_0 satisfies the induction property:

Let $\mathcal{P}(n)$ be a property for natural numbers n ("predicate").

If: (1) $\mathcal{P}(0)$ is true (base case)

(2) $\forall n \in \mathbb{N}_0: \mathcal{P}(n) \rightarrow \mathcal{P}(n+1)$ is true (induction step)

Then: $\mathcal{P}(n)$ is true for all $n \in \mathbb{N}_0$ ($\forall n: \mathcal{P}(n)$ is true)



Proposition: For all $k, m, n \in \mathbb{N}_0$, we have:

$$(k + m) + n = k + (m + n) \quad (\text{associative law})$$

Proof: Use mathematical induction.

$\mathcal{P}(n)$ is given by:

$$\forall k, m \in \mathbb{N}_0: (k + m) + n = k + (m + n)$$

Base case: $\mathcal{P}(0)$ means $\forall k, m \in \mathbb{N}_0 : \underbrace{(k+m)}_{k+m} + 0 = k + \underbrace{(m+0)}_m$

$\Leftrightarrow \forall k, m \in \mathbb{N}_0 : k+m = k+m$ true

Induction step: $(\forall n \in \mathbb{N}_0 : \mathcal{P}(n) \rightarrow \mathcal{P}(n+1))$

Assume $\mathcal{P}(n)$ is true.

$$m + (n+1) = (m+n) + 1 \quad (\text{RD})$$

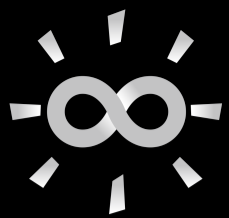
$\mathcal{P}(n+1)$ means $\forall k, m \in \mathbb{N}_0 : (k+m) + (n+1) = k + (m + (n+1))$

Left-hand side: $(k+m) + (n+1) \stackrel{(\text{RD})}{=} ((k+m) + n) + 1$

$\stackrel{\mathcal{P}(n)}{=} (k + (m+n)) + 1$

Right-hand side

$\stackrel{(\text{RD})}{=} k + ((m+n) + 1) \stackrel{(\text{RD})}{=} k + (m + (n+1))$



Start Learning Numbers - Part 4

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

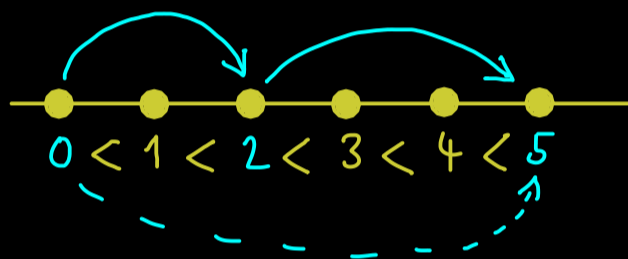
Addition $+$ is a map $\mathbb{N}_0 \times \mathbb{N}_0 \longrightarrow \mathbb{N}_0$ with:

- $m + 0 = m$ (neutral element)
- $(k + m) + n = k + (m + n)$ (associative law)
- $m + n = n + m$ (commutative law)

Ordering:

We write $n \leq m$ if:

$$\exists k \in \mathbb{N}_0 : m = n + k$$



And we write $n < m$ if: $n \leq m \wedge n \neq m$

Properties:

(1) $n \leq n$ (reflexive)

(2) If $n \leq m \wedge m \leq n$, then $n = m$ (antisymmetric)

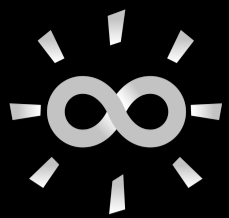
(3) If $n \leq l \wedge l \leq m$, then $n \leq m$ (transitive)

Proof: Assume $n \leq l$ and $l \leq m$ are true. So:

$$\exists k_1 \in \mathbb{N}_0 : l = n + k_1 \quad \text{and} \quad \exists k_2 \in \mathbb{N}_0 : m = l + k_2 \quad \text{are true.}$$

$$\begin{aligned} \text{Therefore:} \quad m &= l + k_2 = (n + k_1) + k_2 \\ &= n + \underbrace{(k_1 + k_2)}_{=: k \in \mathbb{N}_0} = n + k \end{aligned}$$

Therefore: $\exists k \in \mathbb{N}_0 : m = n + k$ is true, so $n \leq m$ is true.



Start Learning Numbers - Part 5

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

$$\underbrace{4 + 4 + 4 + 4 + 4}_{\text{We have 5 of them}} =: 5 \cdot 4$$

We have 5 of them

$$3 + 3 + 3 + 3 + 3 + 3 =: 6 \cdot 3$$

$$4 =: 1 \cdot 4$$

$$0 =: 0 \cdot 4$$

How can we define the multiplication?

Multiplication in \mathbb{N}_0 : map $\mathbb{N}_0 \times \mathbb{N}_0 \longrightarrow \mathbb{N}_0$
 $(n, m) \longmapsto n \cdot m$ defined by

$$0 \cdot m := 0$$

$$(n+1) \cdot m := (n \cdot m) + m$$

(recursive definition)

$$5 \cdot 2 = 2 + 2 + 2 + 2 + 2$$

$$6 \cdot 2 = \underbrace{2 + 2 + 2 + 2 + 2}_{5 \cdot 2} + 2 \quad (\text{Map is well-defined by Dedekind's recursion theorem})$$

Properties: (1) $n \cdot (m \cdot k) = (n \cdot m) \cdot k$ (associative)

(2) $n \cdot m = m \cdot n$ (commutative)

(3) $1 \cdot m = m$ (neutral element)

How to connect + and ·: $n \cdot (m+k) = n \cdot m + n \cdot k$ (distributive)

$$\begin{array}{l} 0 \cdot m := 0 \\ (n+1) \cdot m := (n \cdot m) + m \end{array} \quad (*)$$

Proof by induction: Base case: $n = 0$

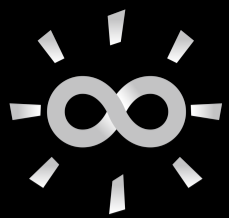
Left-hand side: $0 \cdot (m+k) = 0$ ✓

Right-hand side: $0 \cdot m + 0 \cdot k = 0 + 0 = 0$ ✓

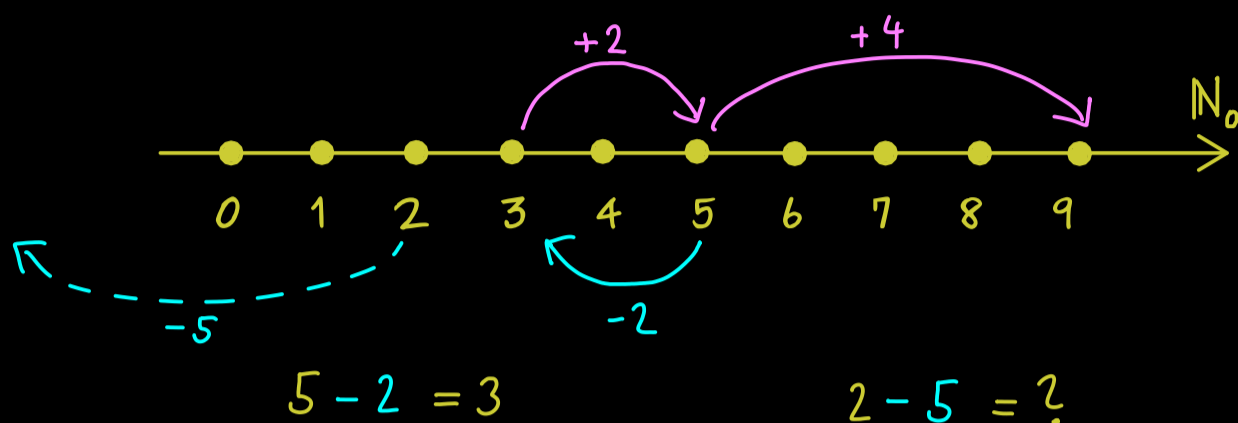
Induction step: Assume $n \cdot (m+k) = n \cdot m + n \cdot k$ holds for n .
(induction hypothesis)

Left-hand side: $(n+1) \cdot (m+k) \stackrel{(*)}{=} n \cdot (m+k) + (m+k)$

$$\begin{aligned} & \stackrel{(i.h.)}{=} n \cdot m + (\underbrace{n \cdot k + m}_{\text{induction hypothesis}}) + k \\ & = (n \cdot m + m) + (n \cdot k + k) \\ & \stackrel{(*)}{=} (n+1) \cdot m + (n+1) \cdot k \quad \leftarrow \text{Right-hand side} \end{aligned}$$



Start Learning Numbers - Part 6



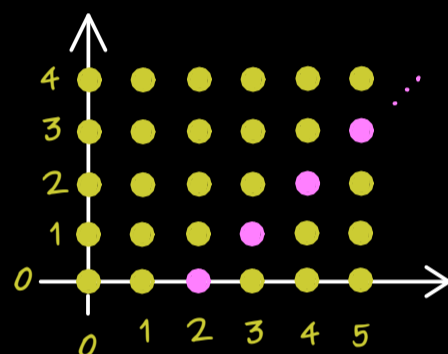
Idea: Look at pairs $(9,5)$, $(5,3)$, $(3,5)$

$$\mathbb{N}_0 \times \mathbb{N}_0 =: \mathbb{N}_0^2$$

$(5,3)$ stands for "5-3"

$(4,2)$ stands for "4-2"

$(5,3) \sim (4,2)$ (equivalent)



\vdots
 $(5,3)$
 $(4,2)$
 $(3,1)$
 $(2,0)$

should be the "same"

$$5-3 = 4-2 \leftarrow \text{not okay}$$

$$5+2 = 4+3 \leftarrow \text{totally okay}$$

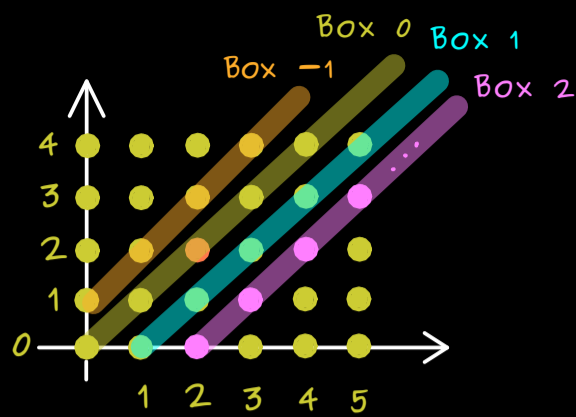
Equivalence relation: We write $(a,b) \sim (x,y)$ if:

$$a+y = x+b$$

Properties: (1) $(a,b) \sim (a,b)$ (reflexive)

(2) If $(a,b) \sim (x,y)$, then $(x,y) \sim (a,b)$. (symmetric)

(3) If $(a,b) \sim (x,y)$ and $(x,y) \sim (c,d)$, then $(a,b) \sim (c,d)$. (transitive)



Property of \mathbb{N}_0 (cancellation):
 If $m + n = \tilde{m} + n$, then $m = \tilde{m}$.

$$\text{Box } 0 = [(2, 2)]_{\sim} := \{(x, y) \in \mathbb{N}_0^2 \mid (x, y) \sim (2, 2)\}$$

is called the equivalence class of $(2, 2)$.

$$\text{Box } 0 = [(0, 0)]_{\sim} = [(2, 2)]_{\sim}$$

$$\text{Box } -1 = [(0, 1)]_{\sim} = [(8, 9)]_{\sim}$$

$$\text{Box } 1 = [(1, 0)]_{\sim} = [(9, 8)]_{\sim}$$

$$\text{Box } -2 = [(0, 2)]_{\sim}$$

$$\vdots$$

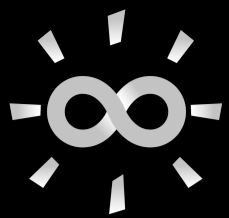
$$\vdots$$

$$\text{Box } 2 = [(2, 0)]_{\sim}$$

$$\vdots$$

$$\vdots$$

$\mathcal{Z} :=$ set of all boxes (equivalence classes)



Start Learning Numbers - Part 7

In \mathbb{N}_0 $4 + x = 0$ is not solvable! No "inverse" of 4.

$$\mathbb{Z} := \{ [(a,b)]_{\sim} \mid (a,b) \in \mathbb{N}_0^2 \} =: \mathbb{N}_0^2 / \sim$$

$$\text{with } [(a,b)]_{\sim} := \{ (x,y) \mid (x,y) \sim (a,b) \}$$

$$\text{and } (x,y) \sim (a,b) \Leftrightarrow x+b = a+y$$

$$[(0,0)]_{\sim} =: 0_{\mathbb{Z}}$$

$$[(0,1)]_{\sim} =: (-1)_{\mathbb{Z}}$$

$$[(1,0)]_{\sim} =: 1_{\mathbb{Z}}$$

$$[(0,2)]_{\sim} =: (-2)_{\mathbb{Z}}$$

$$[(2,0)]_{\sim} =: 2_{\mathbb{Z}}$$

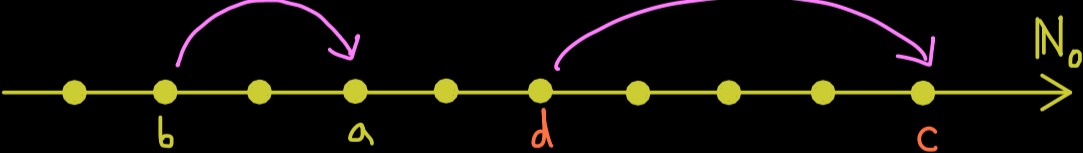
$$\vdots$$
$$\vdots$$

$$\mathbb{Z} = \{ \dots, (-2)_{\mathbb{Z}}, (-1)_{\mathbb{Z}}, 0_{\mathbb{Z}}, 1_{\mathbb{Z}}, 2_{\mathbb{Z}}, \dots \}$$

Question: Is $4_{\mathbb{Z}} + x = 0_{\mathbb{Z}}$ now solvable? And with $x = (-4)_{\mathbb{Z}}$?

First question: How is $+$ as a map $\mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$ defined?

$$[(a,b)]_{\sim} + [(c,d)]_{\sim} := [(a+c, b+d)]_{\sim}$$



well-defined? ✓

Take $(\tilde{a}, \tilde{b}) \sim (a,b)$ and $(\tilde{c}, \tilde{d}) \sim (c,d)$. Then $[(\tilde{a}, \tilde{b})]_{\sim} + [(\tilde{c}, \tilde{d})]_{\sim} = [(\tilde{a} + \tilde{c}, \tilde{b} + \tilde{d})]_{\sim}$

Is $(\tilde{a} + \tilde{c}, \tilde{b} + \tilde{d}) \sim (a+c, b+d)$?

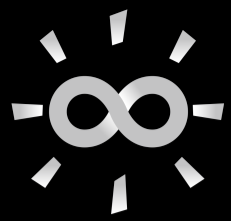
Proof: $(\tilde{a}, \tilde{b}) \sim (a, b) \Leftrightarrow \tilde{a} + b = a + \tilde{b}$
 $(\tilde{c}, \tilde{d}) \sim (c, d) \Leftrightarrow \tilde{c} + d = c + \tilde{d}$ } implies: $\tilde{a} + \tilde{c} + b + d = a + c + \tilde{b} + \tilde{d}$
 $\Leftrightarrow (\tilde{a} + \tilde{c}, \tilde{b} + \tilde{d}) \sim (a + c, b + d)$

Examples: (a) $4_{\mathbb{Z}} + 2_{\mathbb{Z}} = [(4, 0)]_{\sim} + [(2, 0)]_{\sim} = [(6, 0)]_{\sim} = 6_{\mathbb{Z}}$

(b) $4_{\mathbb{Z}} + (-4)_{\mathbb{Z}} = [(4, 0)]_{\sim} + [(0, 4)]_{\sim} = [(4, 4)]_{\sim} = [(0, 0)]_{\sim} = 0_{\mathbb{Z}}$

Properties of \mathbb{Z} together with $+$: $\leftarrow \text{map } \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$

- (a) associative
 - (b) commutative
 - (c) $m + 0_{\mathbb{Z}} = m$ ($0_{\mathbb{Z}}$ is neutral element)
 - (d) For all $m \in \mathbb{Z}$, there is an element $\tilde{m} \in \mathbb{Z}$ with $m + \tilde{m} = 0_{\mathbb{Z}}$
- $\rightarrow (\mathbb{Z}, +)$ is an abelian group



Start Learning Numbers - Part 8

$$\mathbb{Z} = \{ \dots, (-2)_{\mathbb{Z}}, (-1)_{\mathbb{Z}}, 0_{\mathbb{Z}}, 1_{\mathbb{Z}}, 2_{\mathbb{Z}}, \dots \}$$

$$2_{\mathbb{Z}} = [(6,4)]_{\sim} \quad \leftarrow \text{think of "6-4"}$$

$$[(a,b)]_{\sim} \cdot [(c,d)]_{\sim} := [(a \cdot c + b \cdot d, a \cdot d + b \cdot c)]_{\sim} \quad \leftarrow \text{think of "(a-b) \cdot (c-d) = (ac + bd) - (ad + bc)"}$$

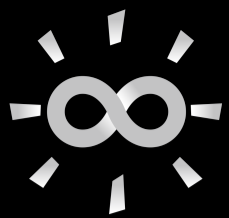
The multiplication is well-defined.

Properties of \mathbb{Z} together with \cdot : $\leftarrow \text{map } \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$

- (a) associative
- (b) commutative
- (c) $1_{\mathbb{Z}} \cdot m = m$ ($1_{\mathbb{Z}}$ is neutral element)
- (d) distributive

Examples: (a) $4_{\mathbb{Z}} \cdot 2_{\mathbb{Z}} = [(4,0)]_{\sim} \cdot [(2,0)]_{\sim} = [(4 \cdot 2 + 0 \cdot 0, 4 \cdot 0 + 0 \cdot 2)]_{\sim} = 8_{\mathbb{Z}}$

(b) $(-4)_{\mathbb{Z}} \cdot (-2)_{\mathbb{Z}} = [(0,4)]_{\sim} \cdot [(0,2)]_{\sim} = [(0 \cdot 0 + 4 \cdot 2, 0 \cdot 2 + 4 \cdot 0)]_{\sim} = 8_{\mathbb{Z}}$



Start Learning Numbers - Part 9

$$\mathbb{Z} = \{\dots -3, -2, -1, 0, 1, 2, 3, 4, \dots\}$$



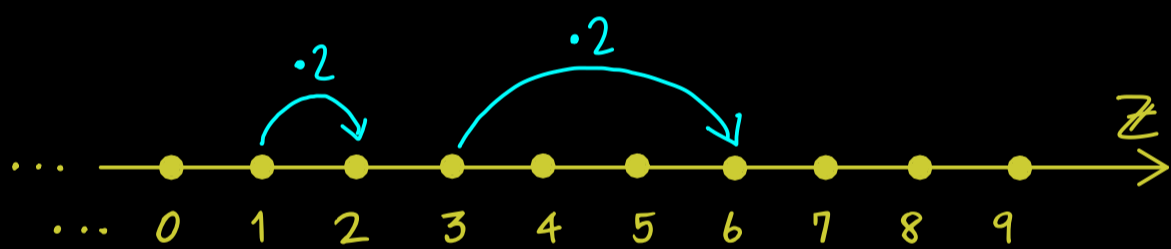
ratio: 3:1 or 3:4 or 1:4

$$\text{fraction: } \frac{3}{4} + \frac{1}{4} = 1$$

solve $4 \cdot x = 1$? \rightsquigarrow We need inverses with respect to \cdot !
Works the same as $(\mathbb{N}_0, +) \rightsquigarrow (\mathbb{Z}, +)$

For $(c, d), (a, b) \in \mathbb{Z} \times \mathbb{Z} \setminus \{0\}$ define:

$$(a, b) \sim (c, d) \quad \text{by} \quad a \cdot d = c \cdot b$$



$$(6, 3) \sim (2, 1) \quad \text{"} \frac{6}{3} = \frac{2}{1} \text{"}$$

$$\mathbb{Q} := (\mathbb{Z} \times \mathbb{Z} \setminus \{0\}) / \sim = \left\{ [(a, b)]_{\sim} \mid (a, b) \in \mathbb{Z} \times \mathbb{Z} \setminus \{0\} \right\} \quad \text{rational numbers}$$

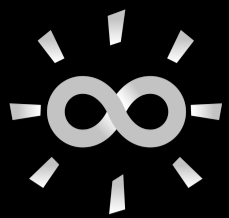
Examples: $[(4, 2)]_{\sim} = [(6, 3)]_{\sim} = [(2, 1)]_{\sim} =: 2_{\mathbb{Q}}$

$$[(-9, -3)]_{\sim} = [(9, 3)]_{\sim} = [(3, 1)]_{\sim} =: 3_{\mathbb{Q}}, \quad [(0, 8)]_{\sim} = [(0, 1)]_{\sim} =: 0_{\mathbb{Q}}$$

$$[(-9, 3)]_{\sim} = [(-3, 1)]_{\sim} =: (-3)_{\mathbb{Q}} \quad \text{We get all integers back!}$$

$$[(2, 8)]_{\sim} = [(1, 4)]_{\sim} =: \left(\frac{1}{4}\right)_{\mathbb{Q}} \rightsquigarrow \text{fractions}$$

Definition: $[(a, b)]_{\sim} =: \frac{a}{b} \quad \left(\frac{2}{8} = \frac{1}{4}\right)$

Start Learning Numbers - Part 10

$$\mathbb{Q} = \left\{ \frac{a}{b} \mid a \in \mathbb{Z}, b \in \mathbb{Z} \setminus \{0\} \right\}, \quad \frac{a}{b} = \frac{c}{d} \Leftrightarrow a \cdot d = c \cdot b$$

Multiplication: $\frac{a}{b} \cdot \frac{c}{d} := \frac{a \cdot c}{b \cdot d}$ well-defined!

For $a \neq 0$, we have: $\frac{a}{b} \cdot \frac{b}{a} = \frac{a \cdot b}{b \cdot a} = \frac{1}{1} (= 1_a)$

solve: $4 \cdot x = 1$? In \mathbb{Q} : $\frac{4}{1} \cdot x = \frac{1}{1}$ is solved by: $x = \frac{1}{4}$

Property: $(\mathbb{Q} \setminus \{0_{\mathbb{Q}}\}, \cdot)$ is an abelian group.

How to define the addition?

We want the distributive law:

$$\boxed{\frac{a}{d} + \frac{c}{d}} = \frac{a}{1} \cdot \frac{1}{d} + \frac{c}{1} \cdot \frac{1}{d} = \left(\frac{a}{1} + \frac{c}{1} \right) \cdot \frac{1}{d} = \boxed{\frac{a+c}{d}}$$

should be defined by:

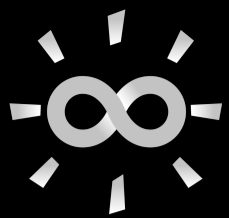
$$\begin{aligned} \frac{a}{b} + \frac{c}{d} &= \frac{a}{b} \cdot \frac{d}{d} + \frac{c}{d} \cdot \frac{b}{b} = \frac{a \cdot d}{1} \cdot \frac{1}{b \cdot d} + \frac{c \cdot b}{1} \cdot \frac{1}{b \cdot d} \\ &= \left(\frac{a \cdot d}{1} + \frac{c \cdot b}{1} \right) \cdot \frac{1}{b \cdot d} = \frac{a \cdot d + c \cdot b}{b \cdot d} \end{aligned}$$

Define: $\frac{a}{b} + \frac{c}{d} := \frac{a \cdot d + c \cdot b}{b \cdot d}$ well-defined!

Proposition: The set \mathbb{Q} together with the operation $+$ and \cdot satisfies:

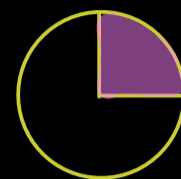
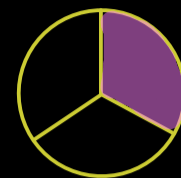
- (1) $(\mathbb{Q}, +)$ is an abelian group
- (2) $(\mathbb{Q} \setminus \{0_{\mathbb{Q}}\}, \cdot)$ is an abelian group
- (3) distributive law

field



Start Learning Numbers - Part 11

$$\mathbb{Q} = \left\{ \frac{a}{b} \mid a \in \mathbb{Z}, b \in \mathbb{Z} \setminus \{0\} \right\}$$

 \leq 

We need to define it

Definition of \leq for \mathbb{Z} : For $a, b \in \mathbb{Z}$, we write $a \leq b$ if $\exists k \in \mathbb{N}_0: a + k = b$

Now: $\frac{1}{4} \leq \frac{1}{3}$ because $3 \leq 4$

Definition of \leq for \mathbb{Q} : For $b > 0$ and $d > 0$

$$\frac{a}{b} \leq \frac{c}{d} \quad \text{defined by} \quad a \cdot d \leq c \cdot b$$

Properties of \leq for \mathbb{Q} : (1) Ordering: reflexive, antisymmetric and transitive.

(2) For all $x, y, z \in \mathbb{Q}$: If $x \leq y$, then $x + z \leq y + z$

(3) For all $x, y, z \in \mathbb{Q}$: If $z \geq 0$ and $x \leq y$, then $x \cdot z \leq y \cdot z$

(4) Total order: For all $x, y \in \mathbb{Q}$, we have $x \leq y$ or $y \leq x$.

(5) Archimedean property: For all $x, \varepsilon \in \mathbb{Q}$ with $x > 0$ and $\varepsilon > 0$,

we have: $n \in \mathbb{N}_0: n \cdot \varepsilon = \varepsilon + \varepsilon + \varepsilon + \dots + \varepsilon > x$

