

The Bright Side of Mathematics

The following pages cover the whole Start Learning Mathematics course of the Bright Side of Mathematics. Please note that the creator lives from generous supporters and would be very happy about a donation. See more here: <https://tbsom.de/support>

Have fun learning mathematics!

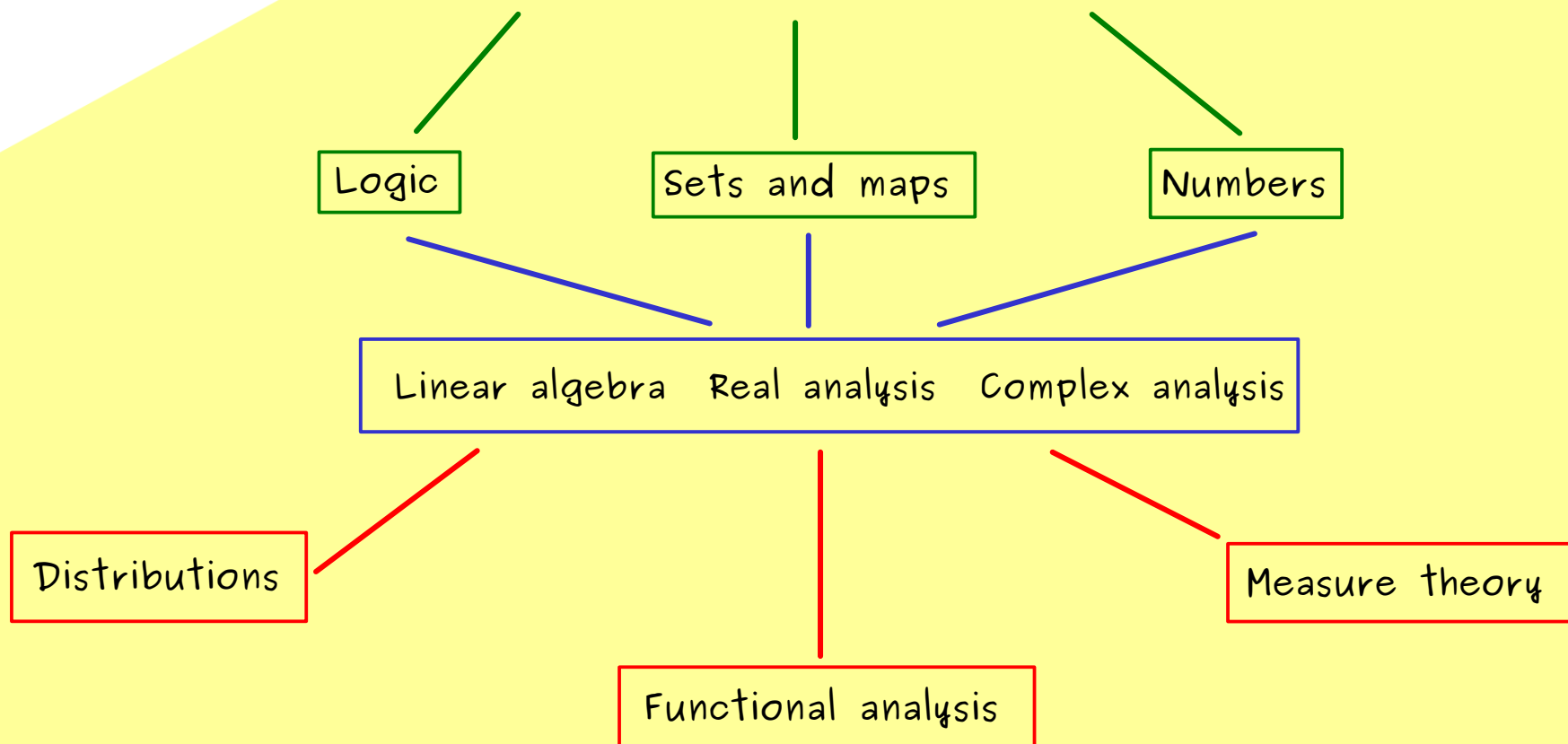
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The Bright Side of Mathematics

Start Learning Mathematics



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Start Learning Logic - Part 1

Logical statement (proposition): Statement that is either **True** or **False**

- Examples:
- (a) Mars is a planet (**True** logical statement)
 - (b) Pluto is a planet (**False** logical statement)
 - (c) $1 + 1 = 2$ (**True** logical statement)
 - (d) The number 5 is smaller than the number 2 (**False** logical statement)
 - (e) Good morning! (**Not** a logical statement)
 - (f) $x + 1 = 1$ (**Not** a logical statement) \rightsquigarrow predicate

Logical operations:

Negation: For a logical statement A ,
 $\neg A$ denotes the negation.

Truth table

A	$\neg A$
T	F
F	T

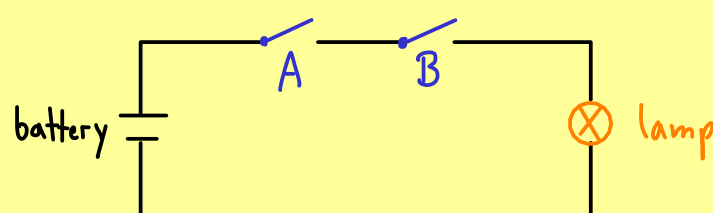
Examples: (a) $A =$ The wine bottle is full
 $\neg A =$ The wine bottle is not full

(b) $A = 2 + 2 = 5$
 $\neg A = 2 + 2 \neq 5$

Conjunction: For two logical statements A, B ,
 $A \wedge B$ denotes the conjunction.

Truth table

A	B	$A \wedge B$
T	T	T
T	F	F
F	T	F
F	F	F





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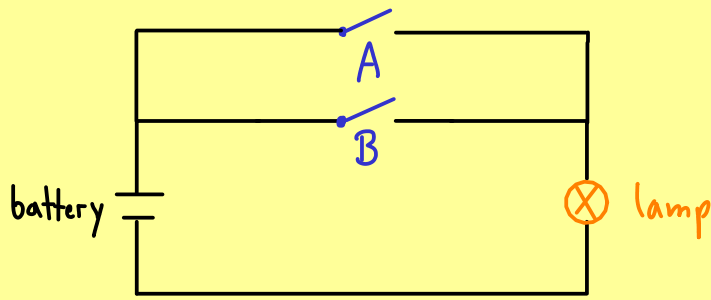
Start Learning Logic - Part 2

Logical statements $A, B \rightsquigarrow$ new logical statements $\neg A, A \vee B$

logical variables A, B
 logical operations \neg, \vee

Logical operations:

Disjunction: For two logical statements A, B , $A \vee B$ denotes the disjunction.



Truth table

A	B	$A \vee B$
T	T	T
T	F	T
F	T	T
F	F	F

Example: $\neg A \vee A$

Truth table

A	$\neg A$	$\neg A \vee A$
T	F	T
F	T	T

We say $\neg A \vee A$ is a tautology.

↳ always true (independent of the truth values of the logical variables that are contained)

Logical equivalence: Two logical statements are called logically equivalent if the truth tables (all possible assignments of truth values for the logical variables) are the same.

Example: $\neg(A \vee B) \Leftrightarrow (\neg A) \wedge (\neg B)$

A	B	$A \vee B$	$\neg A$	$\neg B$	$\neg(A \vee B)$	$\neg A \wedge \neg B$
T	T	T	F	F	F	F
T	F	T	F	T	F	F
F	T	T	T	F	F	F
F	F	F	T	T	T	T



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Start Learning Logic - Part 3

Logical operations:

Conditional: For two logical statements A, B ,
 $A \rightarrow B$ denotes the conditional.

Truth table

A	B	$A \rightarrow B$
T	T	T
T	F	F
F	T	T
F	F	T

\Rightarrow means \rightarrow gives tautology

A	B	$A \wedge B$	$A \wedge B \rightarrow B$
T	T	T	T
T	F	F	T
F	T	F	T
F	F	F	T

We can write:

$$A \wedge B \Rightarrow B$$

Biconditional: For two logical statements A, B ,
 $A \leftrightarrow B$ denotes the biconditional.

Truth table

A	B	$A \leftrightarrow B$
T	T	T
T	F	F
F	T	F
F	F	T

\Leftrightarrow means \leftrightarrow gives tautology

Example: (a) $A \leftrightarrow B \Leftrightarrow (A \rightarrow B) \wedge (B \rightarrow A)$

(b) $A \rightarrow B \Leftrightarrow \neg B \rightarrow \neg A$ (contraposition)

If there is fog, then we have poor visibility

If we don't have poor visibility, there is no fog.

Deduction rules: (how to get new true propositions from other true propositions)

Modus ponens: If $A \rightarrow B$ true and A true, then: B true

A	B	$A \rightarrow B$
T	T	T
T	F	F
F	T	T
F	F	T

Chain syllogism: If $A \rightarrow B$ true and $B \rightarrow C$ true, then: $A \rightarrow C$ true

Reductio ad absurdum: If $A \rightarrow B$ true and $A \rightarrow \neg B$ true, then: $\neg A$ true

The Bright Side of Mathematics

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The Bright Side of Mathematics

Start Learning Sets - Part 1

Propositional Logic
+
Naive set theory

Logic
+
Axioms of
set theory

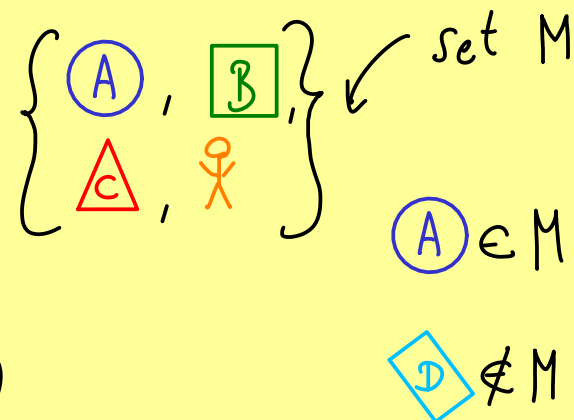
Goal:

doing mathematics

foundation of mathematics

Set: Collection of distinct objects into a whole

Such an object x inside a set M is called an element of M , write: $x \in M$.



If x is not such an object inside the set M , we write: $x \notin M$ means: $\neg(x \in M)$

A set can be defined by giving all its elements:

$$A := \{2, 5, 6\}$$

↑
defined by

Examples: Empty set: $\emptyset := \{\}$

Natural numbers: $\mathbb{N} := \{1, 2, 3, 4, 5, \dots\}$

Natural numbers (including zero): $\mathbb{N}_0 := \{0, 1, 2, 3, 4, \dots\}$

Integers: $\mathbb{Z} := \{\dots, -2, -1, 0, 1, 2, \dots\}$

Rational numbers \mathbb{Q}

Real numbers \mathbb{R}

Complex numbers \mathbb{C}

quantifiers $\forall \exists$ predicates $x \in \mathbb{N}$



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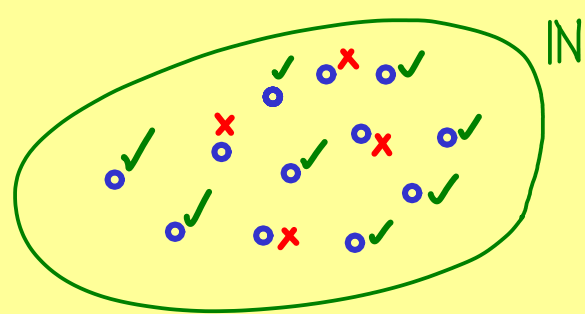
Start Learning Sets - Part 2

- $\boxed{1}$ is an even number false logical statement
 - $\boxed{1}$ is an animal false logical statement
 - $\boxed{1} + 8 = 9$ true logical statement
- } predicates

Predicate: An expression with undetermined variables that ascribes a property to objects filled in for the variables.

Form new sets:

$$\{ x \in \mathbb{N} \mid x \text{ is an even number} \}$$



$$\{ y \in \mathbb{Z} \mid y \in \mathbb{N} \}$$

For $A := \{ \text{Mercury, Venus, Earth, Mars, Jupiter, Saturn, Uranus, Neptune} \}$

form: $\{ p \in A \mid p \text{ has at least 1 confirmed moon} \}$

Quantifiers:

$\forall x$ for all x $\exists x$ it exists x

Predicate: x is a planet

$\forall x : x \text{ is a planet}$ \rightsquigarrow logical statement
false

$\exists x : x \text{ is a planet}$ \rightsquigarrow logical statement
true

Equality for sets: Two sets A, B are the same, written as $A = B$ if

$$\forall x : x \in A \leftrightarrow x \in B \quad \text{is true.}$$

Example: $C := \{ 1, 3, 5 \} = \{ 3, 5, 1 \} =: D$ $1 \in C \leftrightarrow 1 \in D$ true
 $2 \in C \leftrightarrow 2 \in D$ true
 \vdots

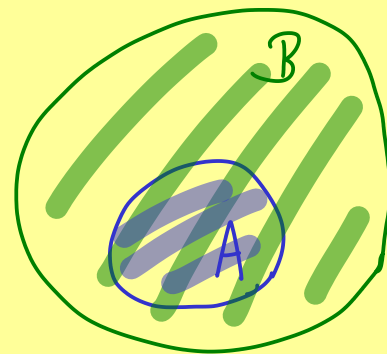
$$\{ 1, 3, 5 \} = \{ 2, 2, 2, 3, 3, 5 \}$$

Subsets: For two sets A, B , we write $A \subseteq B$ if

$$\forall x : x \in A \rightarrow x \in B \quad \text{is true.}$$

short notation: $\forall x \in A : x \in B$

we call A a subset of B . (We can also write $B \supseteq A$)





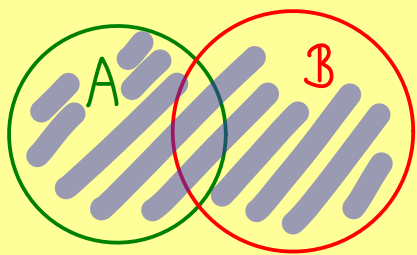
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Start Learning Sets - Part 3

$$A \subseteq B \leftarrow \begin{array}{l} \text{is a superset of } A \\ \text{is a subset of } B \end{array} \rightsquigarrow \begin{array}{l} B \subseteq B \checkmark \\ \emptyset \subseteq B \checkmark \end{array}$$

$$\forall x : x \in \emptyset \rightarrow x \in B$$

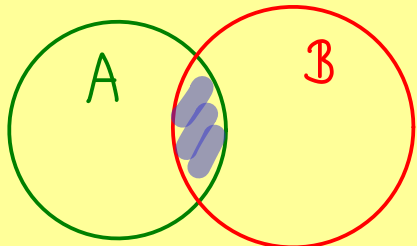
Union:



$$A \cup B := \{x \mid x \in A \vee x \in B\}$$

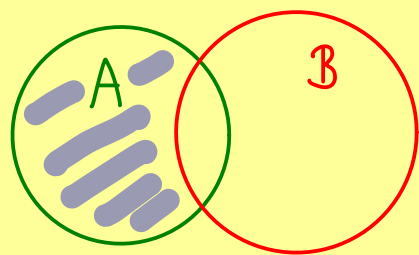
$$(\forall x : x \in A \cup B \leftrightarrow x \in A \vee x \in B) \text{ is true}$$

Intersection:



$$A \cap B := \{x \mid x \in A \wedge x \in B\}$$

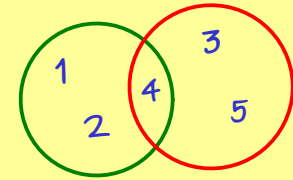
Set difference:



$$A \setminus B := \{x \mid x \in A \wedge x \notin B\}$$

Example:

$$A := \{1, 2, 4\}, \quad B := \{3, 4, 5\}$$

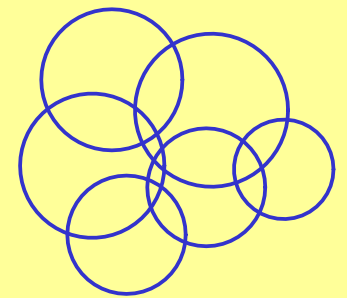


$$A \cup B = \{1, 2, 3, 4, 5\}, \quad A \cap B = \{4\}, \quad A \setminus B = \{1, 2\}$$

Big union:

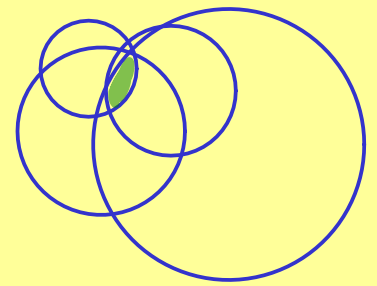
Need: I set, A_i set for each $i \in I$.

$$\bigcup_{i \in I} A_i := \{x \mid \exists i \in I : x \in A_i\}$$



Big intersection:

$$\bigcap_{i \in I} A_i := \{x \mid \forall i \in I : x \in A_i\}$$



Example:

$$A_1 = \{1\}, \quad A_2 = \{2\}, \quad A_3 = \{3\}, \dots$$

$$I = \mathbb{N}, \quad A_i = \{i\}. \text{ Then: } \bigcup_{i \in I} A_i = \{1, 2, 3, \dots\} = \mathbb{N}$$

$$\bigcap_{i \in I} A_i = \emptyset$$

Power set:

$$\text{For a set } A \text{ define } \mathcal{P}(A) := \{X \mid X \subseteq A\}$$

The set of all subsets of A

Example:

$$A = \{1, 2, 3\}, \quad \mathcal{P}(A) = \{\emptyset, \{1, 2, 3\}, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{2, 3\}, \{1, 3\}\}$$

Number of elements:

$$|A| = 3, \quad |\mathcal{P}(A)| = 8 = 2^3$$



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Start Learning Sets - Part 4

Cartesian product: $A \times B$ set of all ordered pairs

$A := \{\Delta, \square, \circ\}$
 $B := \{4, 7\}$

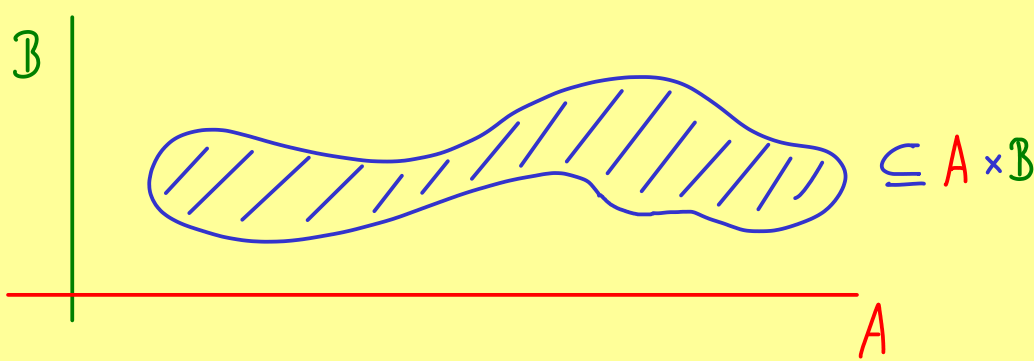
7	$(\Delta, 7)$	$(\square, 7)$	$(\circ, 7)$
4	$(\Delta, 4)$	$(\square, 4)$	$(\circ, 4)$
	Δ	\square	\circ

Definition of ordered pair: For elements x, y write $(x, y) := \{\{x\}, \{x, y\}\}$

$$(x, y) = (\tilde{x}, \tilde{y}) \iff \{x\} = \{\tilde{x}\} \wedge \{x, y\} = \{\tilde{x}, \tilde{y}\}$$

$$\iff x = \tilde{x} \wedge y = \tilde{y}$$

Definition: $A \times B := \{(a, b) \mid a \in A \wedge b \in B\}$



A subset $G_f \subseteq A \times B$ is called a function if

$$(\forall x \forall y \forall \tilde{y} : (x, y) \in G_f \wedge (x, \tilde{y}) \in G_f \rightarrow y = \tilde{y}) \text{ is true.}$$



If also $\forall x \in A : \exists y \in B : (x, y) \in G_f$ is true,

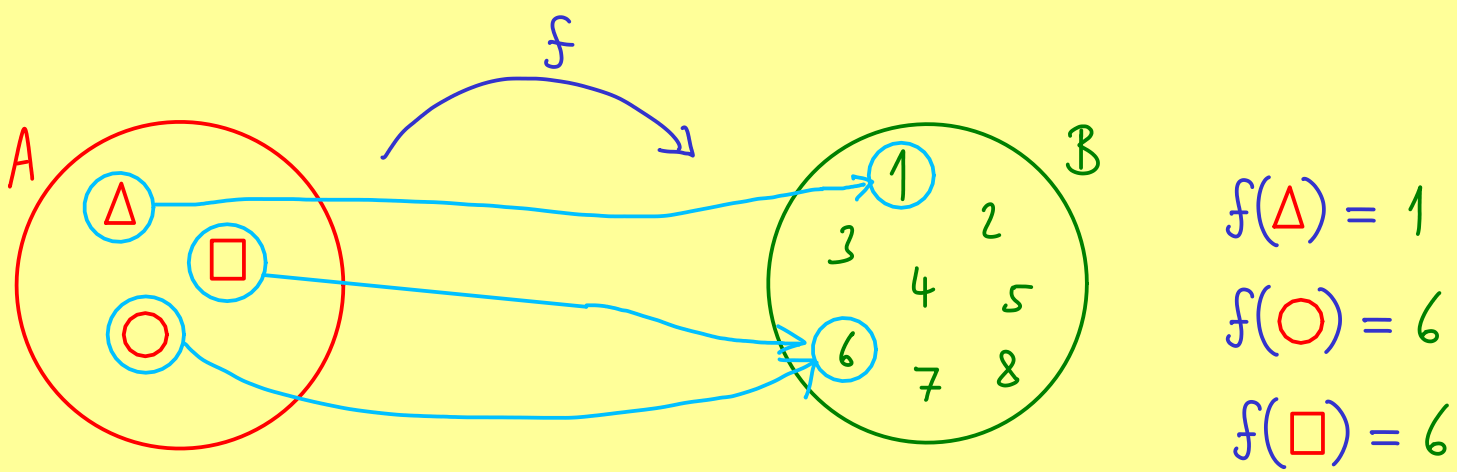
we write:

$f: A \rightarrow B$ and $f(x) = y$ for $(x, y) \in G_f$

↑ domain of f
↑ codomain of f
↑ graph of f

↙ a map from A into B
↘

Example:



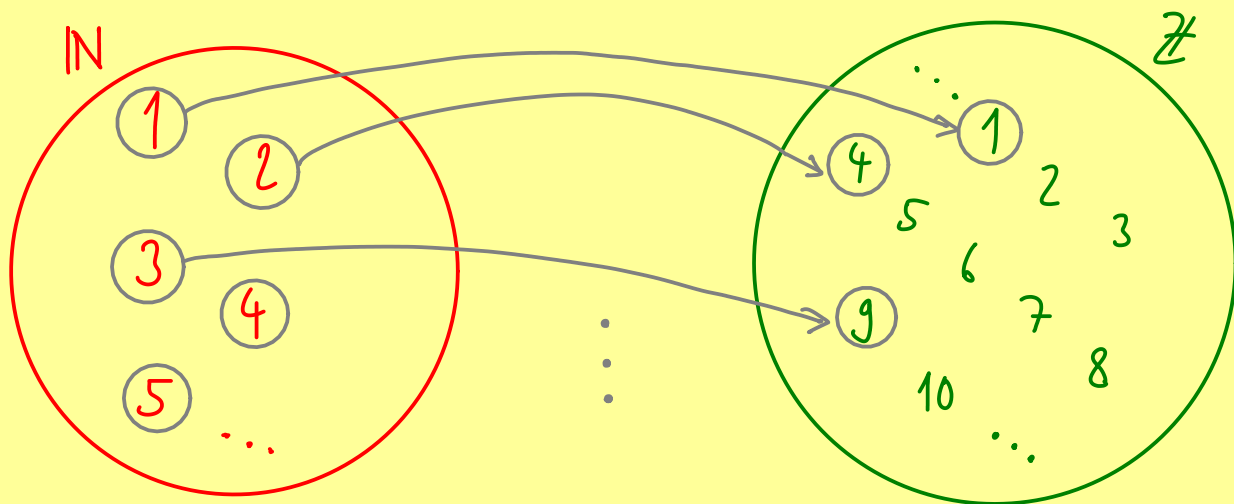


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Start Learning Sets - Part 5

Map: $f: A \rightarrow B$

Example: $f: \mathbb{N} \rightarrow \mathbb{Z}$
 $x \mapsto x^2$ ← new notation for $f(x) = x^2$



Range: $\text{Ran}(f) := \{y \in B \mid \exists x \in A : f(x) = y\}$
 $=: \{f(x) \mid x \in A\}$ (shorter notation)

Example: $f: \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$
 $(x_1, x_2) \mapsto x_1^2 + x_2^2$

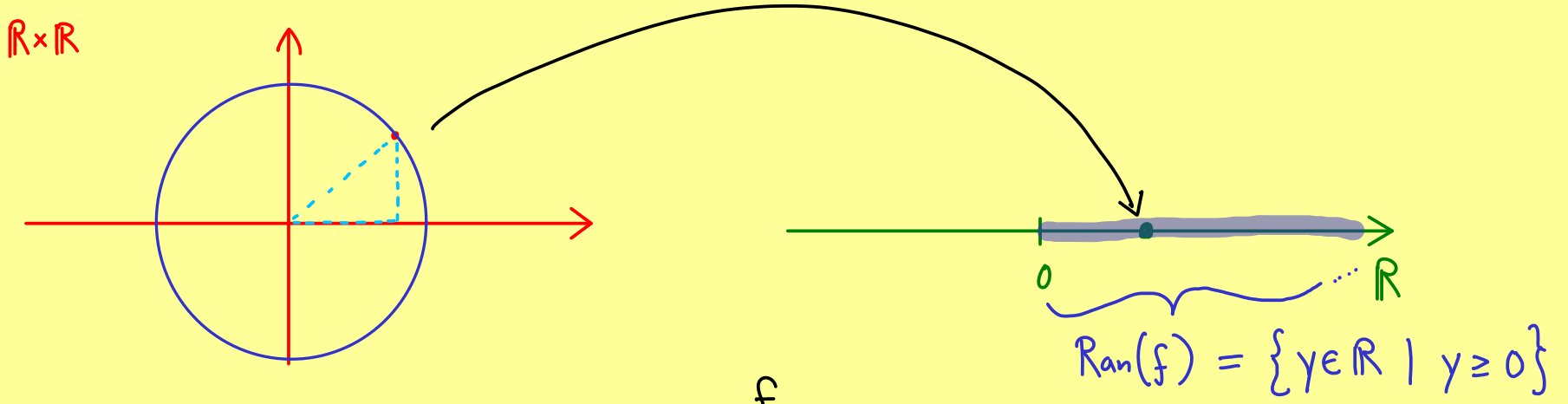
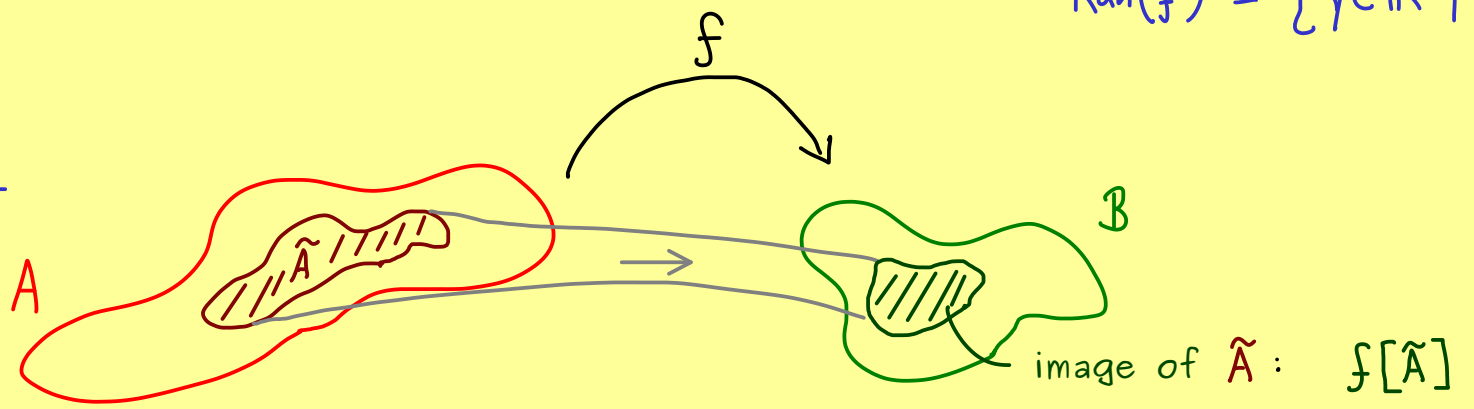


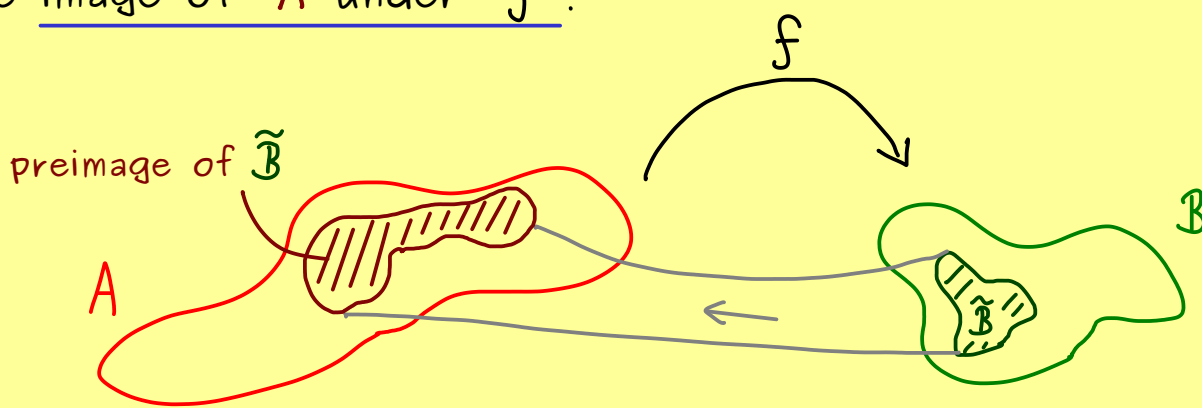
Image and preimage:



For a subset $\tilde{A} \subseteq A$,

$$f[\tilde{A}] := \{y \in B \mid \exists x \in \tilde{A} : f(x) = y\} = \{f(x) \mid x \in \tilde{A}\}$$

denotes the image of \tilde{A} under f .



For $\tilde{B} \subseteq B$,

$$f^{-1}[\tilde{B}] := \{x \in A \mid f(x) \in \tilde{B}\}$$

denotes the preimage of \tilde{B} under f .

Example: $f: \mathbb{N} \rightarrow \mathbb{Z}$

$$x \mapsto \begin{cases} 0 & \text{if } x \text{ even} \\ x & \text{if } x \text{ odd} \end{cases}$$

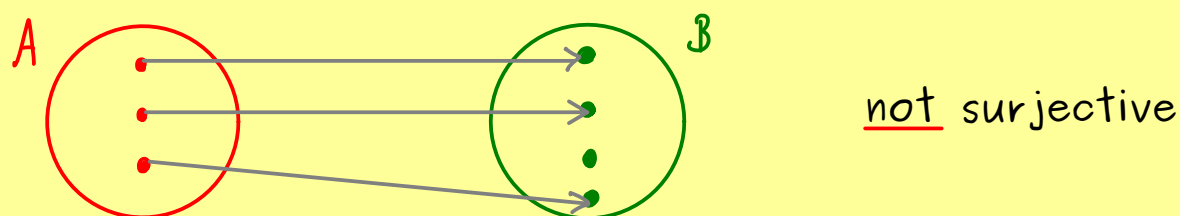
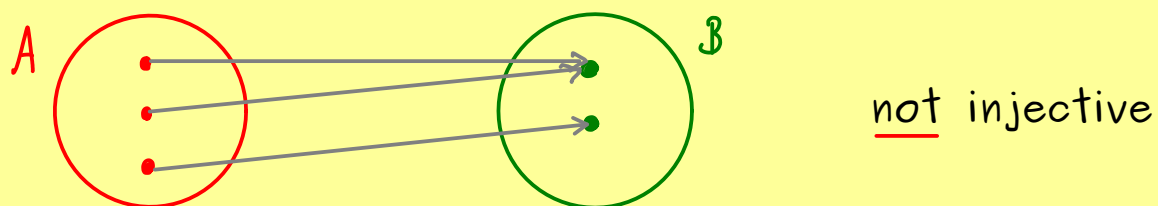
$$f[\{2, 3, 4\}] = \{0, 3\}$$

$$f^{-1}[\{0\}] = \{2, 4, 6, 8, 10, \dots\}$$



The Bright Side of Mathematics

Start Learning Sets - Part 6



Definition: A map $f: A \rightarrow B$ is called:

injective if $\forall x_1, x_2 \in A : (x_1 \neq x_2 \rightarrow f(x_1) \neq f(x_2))$ is true

surjective if $\forall y \in B : \exists x \in A : f(x) = y$ is true

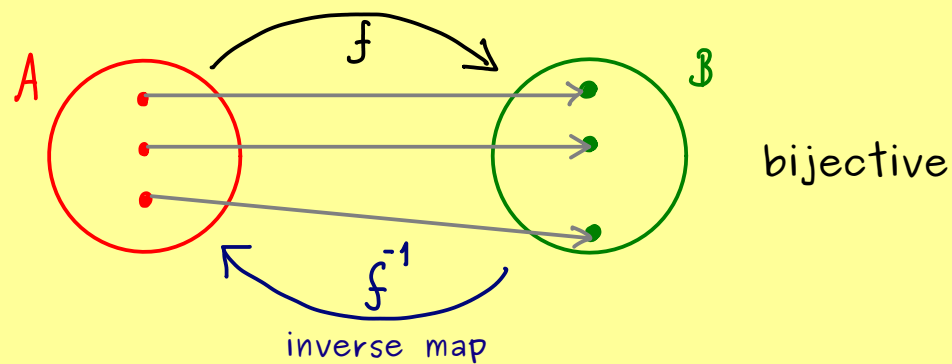
Remember:

surjective: Each $y \in B$ gets at least one arrow.

injective: Each $y \in B$ gets at most one arrow.

injective + surjective Each $y \in B$ gets exactly one arrow.

bijjective (1:1)
 \Leftrightarrow
 invertible



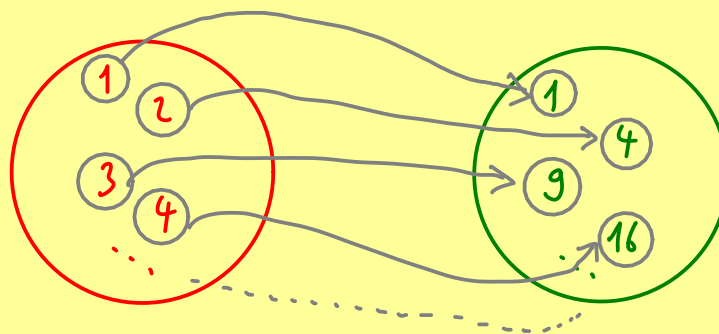
$$f^{-1}: B \rightarrow A,$$

$$f^{-1}(y) := x \quad \text{if} \quad f(x) = y$$

Example:

$$f: \mathbb{N} \rightarrow \{1, 4, 9, 16, 25, 36, \dots\}$$

$$x \mapsto x^2$$



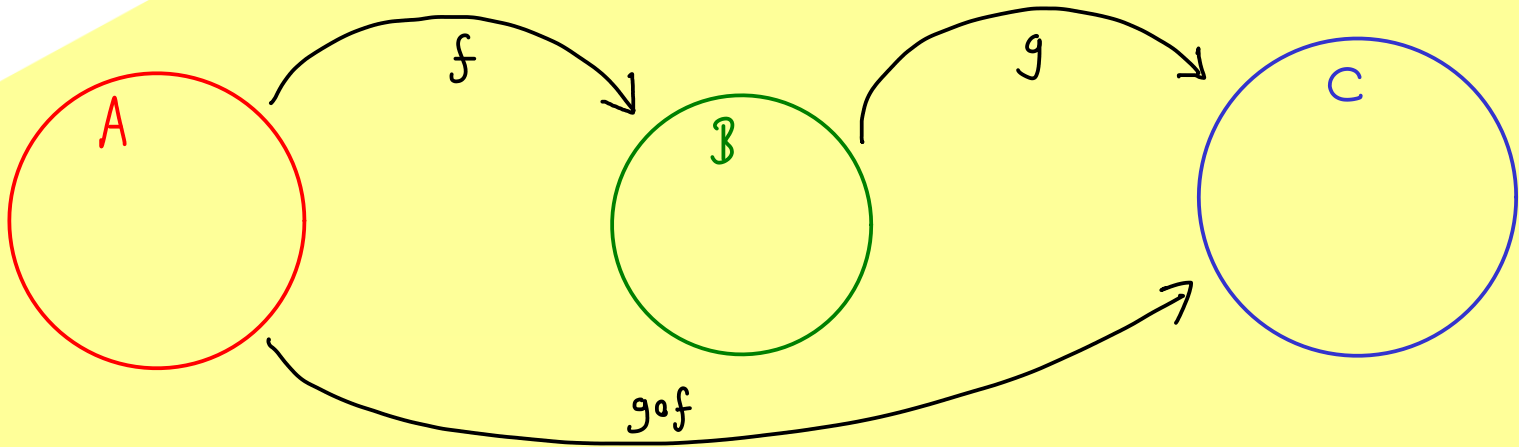
$$f^{-1}: \{1, 4, 9, 16, 25, 36, \dots\} \rightarrow \mathbb{N}$$

$$y \mapsto \sqrt{y}$$



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Start Learning Sets - Part 7



For $f: A \rightarrow B$ and $g: B \rightarrow C$ define:

$$\left. \begin{aligned} g \circ f: A &\rightarrow C \\ x &\mapsto g(f(x)) \end{aligned} \right\} \text{ called the } \underline{\text{composition}} \text{ } g \text{ with } f$$

Examples:

(1)

$(g \circ f)(1) = 9$
 $(g \circ f)(3) = 9$

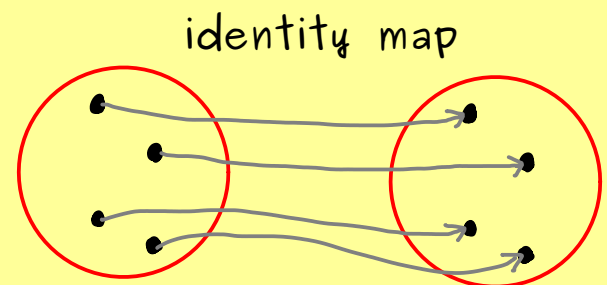
$(g \circ f)(2) = 9$
 $(g \circ f)(4) = 9$

$g(8) = 9$
 $g(2) = 2$
 $f(1) = 8$
 $f(2) = 8$
 $f(3) = 8$
 $f(4) = 8$

(2) $f: \mathbb{R} \rightarrow \mathbb{R}, \quad g: \mathbb{R} \rightarrow \mathbb{R}$
 $x \mapsto x^2 \qquad \qquad x \mapsto \sin(x)$

$\rightsquigarrow (g \circ f)(x) = \sin(x^2) \quad \text{and} \quad (f \circ g)(x) = (\sin(x))^2$

For any set A , we define: $id_A: A \rightarrow A$
 $x \mapsto x$



For $f: A \rightarrow B$ bijective, we have:

$$\begin{aligned} f \circ f^{-1} &= id_B \\ f^{-1} \circ f &= id_A \end{aligned}$$

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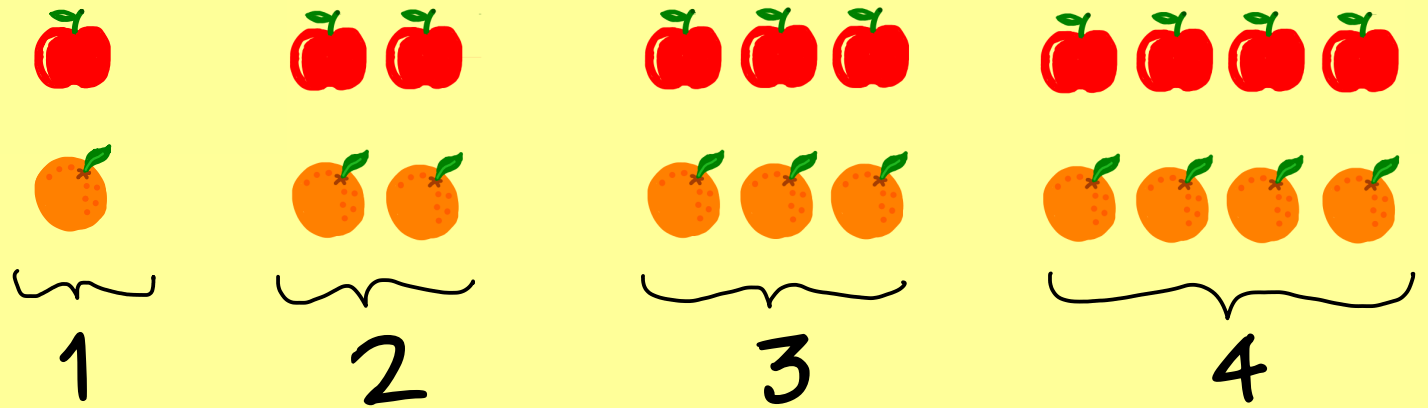
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The Bright Side of Mathematics

Start Learning Numbers - Part 1

Natural numbers



$$\mathbb{N} = \{1, 2, 3, 4, \dots\}$$

$$\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$$

$$0 := \emptyset \text{ empty set}$$

$$1 := \{0\} \text{ set with one element}$$

$$2 := \{0, 1\} \text{ set with two elements}$$

$$3 := \{0, 1, 2\} \text{ set with three elements}$$

$$4 := \{0, 1, 2, 3\} = 3 \cup \{3\}$$

⋮

Axiom: There is a set \mathbb{N}_0 with the properties: (a) $0 \in \mathbb{N}_0$

(b) $\forall x: x \in \mathbb{N}_0 \rightarrow x \cup \{x\} \in \mathbb{N}_0$

And \mathbb{N}_0 is the smallest set having these two properties.

Successor map: $s: \mathbb{N}_0 \rightarrow \mathbb{N}_0$
 $x \mapsto x \cup \{x\}$, $s(6) = 7$



The Bright Side of Mathematics

Start Learning Numbers - Part 2

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

Properties of \mathbb{N}_0 : (1) $0 \in \mathbb{N}_0$

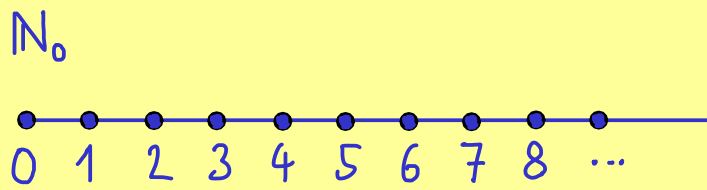
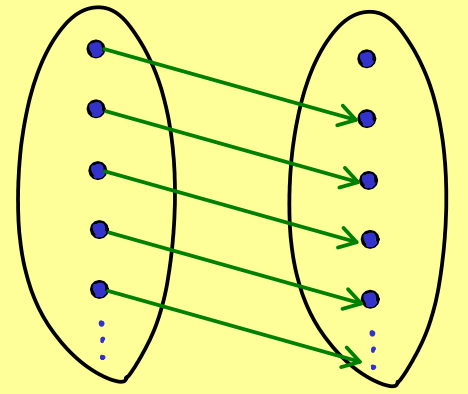
(2) There is a map $s: \mathbb{N}_0 \rightarrow \mathbb{N}_0$ that satisfies:

(2a) s is injective

(2b) $0 \notin \text{Ran}(s) = s[\mathbb{N}_0]$

(2c) If $M \subseteq \mathbb{N}_0$ with
 $0 \in M$ and $s[M] \subseteq M$,
 then $M = \mathbb{N}_0$.

(mathematical induction)



Addition in \mathbb{N}_0 : map $\mathbb{N}_0 \times \mathbb{N}_0 \rightarrow \mathbb{N}_0$
 $(m, n) \mapsto m + n$

How is it defined? $2 + 4 := 6$

$$\boxed{m + 0 := m}, \quad m + 1 := s(m), \quad m + 2 := s(m + 1)$$

Recursive definition:

$$\boxed{m + s(n) := s(m + n)}$$

$$2 + 5 = 2 + s(4) = s(2 + 4) = s(6) = 7$$

Dedekind's principle of recursive definition:

For a set A , $a \in A$ and $h: A \rightarrow A$, then there exists a unique map
 $f: \mathbb{N}_0 \rightarrow A$ with $f(0) = a$ and $f(s(n)) = h(f(n))$.

(" $a, h(a), h(h(a)), h(h(h(a))), \dots$ ")



The Bright Side of Mathematics

Start Learning Numbers - Part 3

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

Each $n \in \mathbb{N}_0$ has a unique successor:

$$s: \mathbb{N}_0 \rightarrow \mathbb{N}_0, \quad s(n) = n + 1$$

We already know: $m + (n + 1) = (m + n) + 1$ (RD)

Mathematical induction:

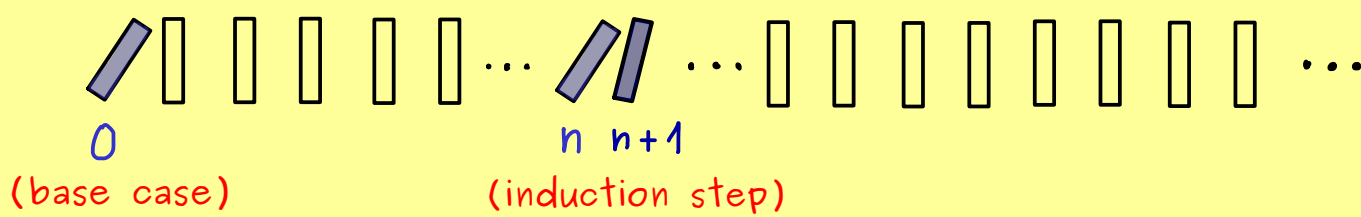
\mathbb{N}_0 satisfies the induction property:

Let $\mathcal{P}(n)$ be a property for natural numbers n ("predicate").

If: (1) $\mathcal{P}(0)$ is true (base case)

(2) $\forall n \in \mathbb{N}_0: \mathcal{P}(n) \rightarrow \mathcal{P}(n+1)$ is true (induction step)

Then: $\mathcal{P}(n)$ is true for all $n \in \mathbb{N}_0$ ($\forall n: \mathcal{P}(n)$ is true)



Proposition: For all $k, m, n \in \mathbb{N}_0$, we have:

$$(k + m) + n = k + (m + n) \quad (\text{associative law})$$

Proof: Use mathematical induction.

$\mathcal{P}(n)$ is given by:

$$\forall k, m \in \mathbb{N}_0: (k + m) + n = k + (m + n)$$

Base case: $\mathcal{P}(0)$ means $\forall k, m \in \mathbb{N}_0: \underbrace{(k + m) + 0}_{k + m} = k + \underbrace{(m + 0)}_m$

$$\Leftrightarrow \forall k, m \in \mathbb{N}_0: k + m = k + m \quad \underline{\text{true}}$$

Induction step: ($\forall n \in \mathbb{N}_0: \mathcal{P}(n) \rightarrow \mathcal{P}(n+1)$)

Assume $\mathcal{P}(n)$ is true.

$$m + (n + 1) = (m + n) + 1 \quad (\text{RD})$$

$$\mathcal{P}(n+1) \text{ means } \forall k, m \in \mathbb{N}_0: (k + m) + (n + 1) = k + (m + (n + 1))$$

$$\text{Left-hand side: } (k + m) + (n + 1) \stackrel{(\text{RD})}{=} ((k + m) + n) + 1$$

$$\stackrel{\mathcal{P}(n)}{=} (k + (m + n)) + 1 \quad \text{Right-hand side}$$

$$\stackrel{(\text{RD})}{=} k + ((m + n) + 1) \stackrel{(\text{RD})}{=} k + (m + (n + 1))$$



The Bright Side of Mathematics

Start Learning Numbers - Part 4

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

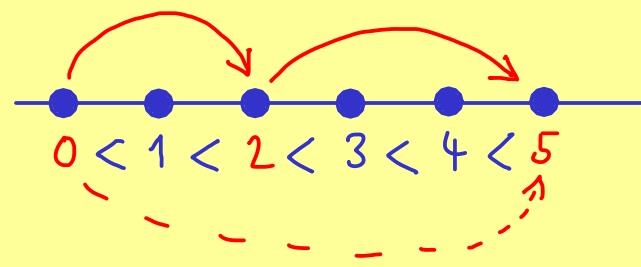
Addition $+$ is a map $\mathbb{N}_0 \times \mathbb{N}_0 \rightarrow \mathbb{N}_0$ with:

- $m + 0 = m$ (neutral element)
- $(k+m) + n = k + (m+n)$ (associative law)
- $m+n = n+m$ (commutative law)

Ordering:

We write $n \leq m$ if:

$$\exists k \in \mathbb{N}_0 : m = n + k$$



And we write $n < m$ if: $n \leq m \wedge n \neq m$

Properties:

- (1) $n \leq n$ (reflexive)
- (2) If $n \leq m \wedge m \leq n$, then $n = m$ (antisymmetric)
- (3) If $n \leq l \wedge l \leq m$, then $n \leq m$ (transitive)

Proof: Assume $n \leq l$ and $l \leq m$ are true. So:

$$\exists k_1 \in \mathbb{N}_0 : l = n + k_1 \quad \text{and} \quad \exists k_2 \in \mathbb{N}_0 : m = l + k_2 \quad \text{are true.}$$

$$\begin{aligned} \text{Therefore: } m &= l + k_2 = (n + k_1) + k_2 \\ &= n + \underbrace{(k_1 + k_2)}_{=: k \in \mathbb{N}_0} = n + k \end{aligned}$$

Therefore: $\exists k \in \mathbb{N}_0 : m = n + k$ is true, so $n \leq m$ is true.



The Bright Side of Mathematics

Start Learning Numbers - Part 5

Natural numbers: $\mathbb{N}_0 = \{0, 1, 2, 3, 4, \dots\}$

$$\underbrace{4 + 4 + 4 + 4 + 4}_{\text{We have 5 of them}} =: 5 \cdot 4$$

We have 5 of them

$$3 + 3 + 3 + 3 + 3 + 3 =: 6 \cdot 3$$

$$4 =: 1 \cdot 4$$

$$0 =: 0 \cdot 4$$

How can we define the multiplication?

Multiplication in \mathbb{N}_0 : map $\mathbb{N}_0 \times \mathbb{N}_0 \rightarrow \mathbb{N}_0$
 $(n, m) \mapsto n \cdot m$ defined by

$$0 \cdot m := 0$$

$$(n+1) \cdot m := (n \cdot m) + m$$

(recursive definition)

$$5 \cdot 2 = 2 + 2 + 2 + 2 + 2$$

$$\text{successor} \quad 6 \cdot 2 = \underbrace{2 + 2 + 2 + 2 + 2}_{5 \cdot 2} + 2 \quad (\text{Map is well-defined by Dedekind's recursion theorem})$$

- Properties:
- (1) $n \cdot (m \cdot k) = (n \cdot m) \cdot k$ (associative)
 - (2) $n \cdot m = m \cdot n$ (commutative)
 - (3) $1 \cdot m = m$ (neutral element)

How to connect + and \cdot : $n \cdot (m+k) = n \cdot m + n \cdot k$ (distributive)

$$0 \cdot m := 0$$

$$(n+1) \cdot m := (n \cdot m) + m$$

(*)

Proof by induction: Base case: $n = 0$

$$\text{Left-hand side: } 0 \cdot (m+k) = 0 \quad \checkmark$$

$$\text{Right-hand side: } 0 \cdot m + 0 \cdot k = 0 + 0 = 0 \quad \checkmark$$

Induction step: Assume $n \cdot (m+k) = n \cdot m + n \cdot k$ holds for n .
 (induction hypothesis)

$$\text{Left-hand side: } (n+1) \cdot (m+k) \stackrel{(*)}{=} n \cdot (m+k) + (m+k)$$

$$\stackrel{\text{(i.h.)}}{=} n \cdot m + (\underbrace{n \cdot k + m}_{\text{induction hypothesis}}) + k$$

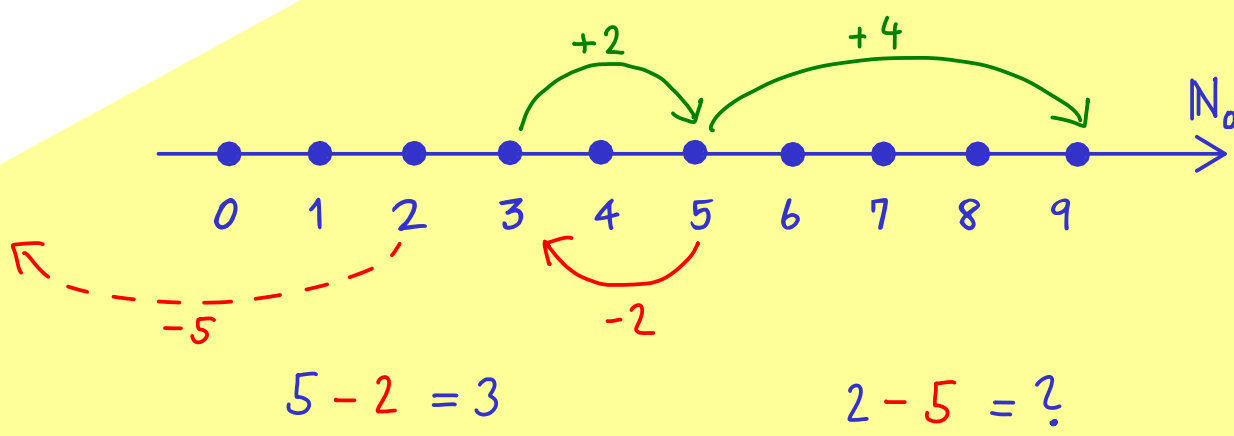
$$= (n \cdot m + m) + (n \cdot k + k)$$

$$\stackrel{(*)}{=} (n+1) \cdot m + (n+1) \cdot k \quad \leftarrow \text{Right-hand side}$$



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Start Learning Numbers - Part 6



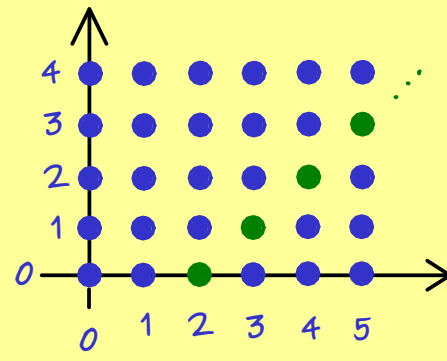
Idea: Look at pairs $(9,5)$, $(5,3)$, $(3,5)$

$$\mathbb{N}_0 \times \mathbb{N}_0 =: \mathbb{N}_0^2$$

$(5,3)$ stands for "5-3"

$(4,2)$ stands for "4-2"

$(5,3) \sim (4,2)$ (equivalent)



\vdots
 $(5,3)$
 $(4,2)$
 $(3,1)$
 $(2,0)$
 should be the "same"

$5-3 = 4-2$ ← not okay

$5+2 = 4+3$ ← totally okay

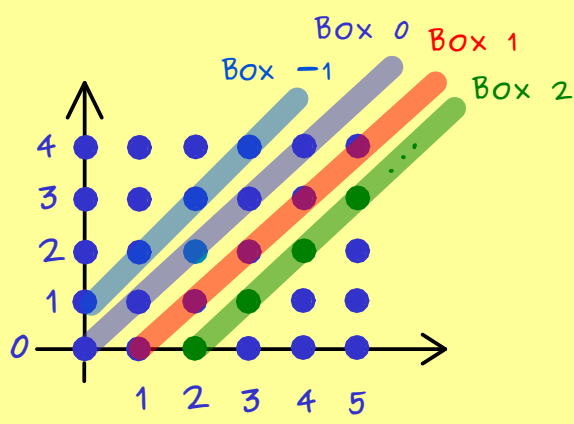
Equivalence relation: We write $(a,b) \sim (x,y)$ if:

$$a+y = x+b$$

Properties: (1) $(a,b) \sim (a,b)$ (reflexive)

(2) If $(a,b) \sim (x,y)$, then $(x,y) \sim (a,b)$. (symmetric)

(3) If $(a,b) \sim (x,y)$ and $(x,y) \sim (c,d)$, then $(a,b) \sim (c,d)$. (transitive)



Property of \mathbb{N}_0 (cancellation):
 If $m+n = \tilde{m}+n$, then $m = \tilde{m}$.

$$\text{Box } 0 = [(2,2)]_{\sim} := \{(x,y) \in \mathbb{N}_0^2 \mid (x,y) \sim (2,2)\}$$

is called the equivalence class of $(2,2)$.

$$\text{Box } 0 = [(0,0)]_{\sim} = [(2,2)]_{\sim}$$

$$\text{Box } -1 = [(0,1)]_{\sim} = [(8,9)]_{\sim}$$

$$\text{Box } 1 = [(1,0)]_{\sim} = [(9,8)]_{\sim}$$

$$\text{Box } -2 = [(0,2)]_{\sim}$$

$$\text{Box } 2 = [(2,0)]_{\sim}$$

$\mathbb{Z} :=$ set of all boxes (equivalence classes)



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Start Learning Numbers - Part 7

In \mathbb{N}_0 $4 + x = 0$ is not solvable! No "inverse" of 4.

$$\mathbb{Z} := \{ [(a,b)]_{\sim} \mid (a,b) \in \mathbb{N}_0^2 \} =: \mathbb{N}_0^2 / \sim$$

$$\text{with } [(a,b)]_{\sim} := \{ (x,y) \mid (x,y) \sim (a,b) \}$$

$$\text{and } (x,y) \sim (a,b) \Leftrightarrow x+b = a+y$$

$$[(0,0)]_{\sim} =: 0_{\mathbb{Z}}$$

$$[(0,1)]_{\sim} =: (-1)_{\mathbb{Z}}$$

$$[(1,0)]_{\sim} =: 1_{\mathbb{Z}}$$

$$[(0,2)]_{\sim} =: (-2)_{\mathbb{Z}}$$

$$[(2,0)]_{\sim} =: 2_{\mathbb{Z}}$$

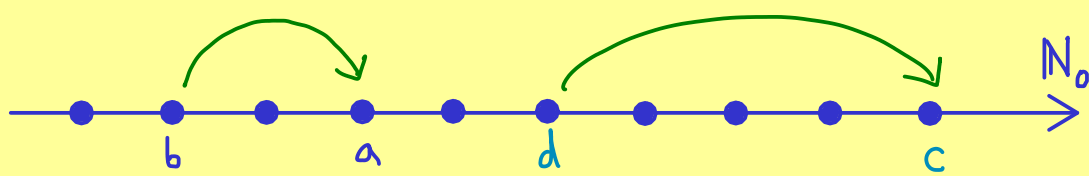
$$\vdots$$

$$\mathbb{Z} = \{ \dots, (-2)_{\mathbb{Z}}, (-1)_{\mathbb{Z}}, 0_{\mathbb{Z}}, 1_{\mathbb{Z}}, 2_{\mathbb{Z}}, \dots \}$$

Question: Is $4_{\mathbb{Z}} + x = 0_{\mathbb{Z}}$ now solvable? And with $x = (-4)_{\mathbb{Z}}$?

First question: How is $+$ as a map $\mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$ defined?

$$[(a,b)]_{\sim} + [(c,d)]_{\sim} := [(a+c, b+d)]_{\sim}$$



well-defined? ✓

Take $(\tilde{a}, \tilde{b}) \sim (a,b)$ and $(\tilde{c}, \tilde{d}) \sim (c,d)$. Then $[(\tilde{a}, \tilde{b})]_{\sim} + [(\tilde{c}, \tilde{d})]_{\sim} = [(\tilde{a} + \tilde{c}, \tilde{b} + \tilde{d})]_{\sim}$

Is $(\tilde{a} + \tilde{c}, \tilde{b} + \tilde{d}) \sim (a+c, b+d)$?

$$\begin{array}{l} \text{Proof: } (\tilde{a}, \tilde{b}) \sim (a,b) \Leftrightarrow \tilde{a} + b = a + \tilde{b} \\ (\tilde{c}, \tilde{d}) \sim (c,d) \Leftrightarrow \tilde{c} + d = c + \tilde{d} \end{array} \left. \vphantom{\begin{array}{l} \text{Proof: } (\tilde{a}, \tilde{b}) \sim (a,b) \Leftrightarrow \tilde{a} + b = a + \tilde{b} \\ (\tilde{c}, \tilde{d}) \sim (c,d) \Leftrightarrow \tilde{c} + d = c + \tilde{d} \end{array}} \right\} \text{implies: } \tilde{a} + \tilde{c} + b + d = a + c + \tilde{b} + \tilde{d} \\ \Leftrightarrow (\tilde{a} + \tilde{c}, \tilde{b} + \tilde{d}) \sim (a+c, b+d)$$

Examples: (a) $4_{\mathbb{Z}} + 2_{\mathbb{Z}} = [(4,0)]_{\sim} + [(2,0)]_{\sim} = [(6,0)]_{\sim} = 6_{\mathbb{Z}}$

(b) $4_{\mathbb{Z}} + (-4)_{\mathbb{Z}} = [(4,0)]_{\sim} + [(0,4)]_{\sim} = [(4,4)]_{\sim} = [(0,0)]_{\sim} = 0_{\mathbb{Z}}$

Properties of \mathbb{Z} together with $+$: \leftarrow map $\mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$

(a) associative

(b) commutative

(c) $m + 0_{\mathbb{Z}} = m$ ($0_{\mathbb{Z}}$ is neutral element)

(d) For all $m \in \mathbb{Z}$, there is an element $\tilde{m} \in \mathbb{Z}$ with $m + \tilde{m} = 0_{\mathbb{Z}}$

$(\mathbb{Z}, +)$ is an abelian group



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Start Learning Numbers - Part 8

$$\mathbb{Z} = \{ \dots, (-2)_{\mathbb{Z}}, (-1)_{\mathbb{Z}}, 0_{\mathbb{Z}}, 1_{\mathbb{Z}}, 2_{\mathbb{Z}}, \dots \}$$

$$2_{\mathbb{Z}} = [(6,4)]_{\sim} \quad \leftarrow \text{think of "6-4"}$$

$$[(a,b)]_{\sim} \cdot [(c,d)]_{\sim} := [(a \cdot c + b \cdot d, a \cdot d + b \cdot c)]_{\sim} \quad \leftarrow \text{think of "(a-b) \cdot (c-d) = (ac + bd) - (ad + bc)"}$$

The multiplication is well-defined.

$$\leftarrow \text{map } \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$$

Properties of \mathbb{Z} together with \cdot :

- (a) associative
- (b) commutative
- (c) $1_{\mathbb{Z}} \cdot m = m$ ($1_{\mathbb{Z}}$ is neutral element)
- (d) distributive

Examples: (a) $4_{\mathbb{Z}} \cdot 2_{\mathbb{Z}} = [(4,0)]_{\sim} \cdot [(2,0)]_{\sim} = [(4 \cdot 2 + 0 \cdot 0, 4 \cdot 0 + 0 \cdot 2)]_{\sim} = 8_{\mathbb{Z}}$

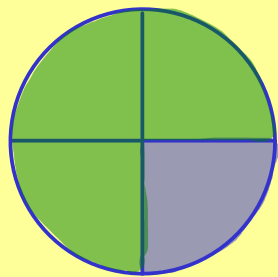
(b) $(-4)_{\mathbb{Z}} \cdot (-2)_{\mathbb{Z}} = [(0,4)]_{\sim} \cdot [(0,2)]_{\sim} = [(0 \cdot 0 + 4 \cdot 2, 0 \cdot 2 + 4 \cdot 0)]_{\sim} = 8_{\mathbb{Z}}$



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Start Learning Numbers - Part 9

$$\mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, 3, 4, \dots\}$$



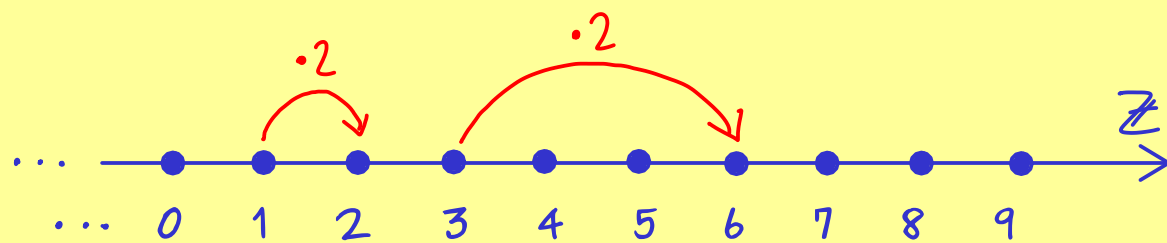
ratio: 3:1 or 3:4 or 1:4

$$\text{fraction: } \frac{3}{4} + \frac{1}{4} = 1$$

solve $4 \cdot x = 1$? \rightsquigarrow We need inverses with respect to \cdot !
Works the same as $(\mathbb{N}_0, +) \rightsquigarrow (\mathbb{Z}, +)$

For $(c, d), (a, b) \in \mathbb{Z} \times \mathbb{Z} \setminus \{0\}$ define:

$$(a, b) \sim (c, d) \quad \text{by} \quad a \cdot d = c \cdot b$$



$$(6, 3) \sim (2, 1) \quad \text{"} \frac{6}{3} = \frac{2}{1} \text{"}$$

$$\mathbb{Q} := (\mathbb{Z} \times \mathbb{Z} \setminus \{0\}) / \sim = \left\{ [(a, b)]_{\sim} \mid (a, b) \in \mathbb{Z} \times \mathbb{Z} \setminus \{0\} \right\} \quad \text{rational numbers}$$

Examples: $[(4, 2)]_{\sim} = [(6, 3)]_{\sim} = [(2, 1)]_{\sim} =: 2_{\mathbb{Q}}$

$$[(-9, -3)]_{\sim} = [(9, 3)]_{\sim} = [(3, 1)]_{\sim} =: 3_{\mathbb{Q}}, \quad [(0, 8)]_{\sim} = [(0, 1)]_{\sim} =: 0_{\mathbb{Q}}$$

$$[(-9, 3)]_{\sim} = [(-3, 1)]_{\sim} =: (-3)_{\mathbb{Q}}$$

We get all integers back!

$$[(2, 8)]_{\sim} = [(1, 4)]_{\sim} =: \left(\frac{1}{4}\right)_{\mathbb{Q}} \rightsquigarrow \text{fractions}$$

Definition: $[(a, b)]_{\sim} =: \frac{a}{b} \quad \left(\frac{2}{8} = \frac{1}{4} \right)$



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Start Learning Numbers - Part 10

$$\mathbb{Q} = \left\{ \frac{a}{b} \mid a \in \mathbb{Z}, b \in \mathbb{Z} \setminus \{0\} \right\}, \quad \frac{a}{b} = \frac{c}{d} \Leftrightarrow a \cdot d = c \cdot b$$

Multiplication: $\frac{a}{b} \cdot \frac{c}{d} := \frac{a \cdot c}{b \cdot d}$ well-defined:

For $a \neq 0$, we have: $\frac{a}{b} \cdot \frac{b}{a} = \frac{a \cdot b}{b \cdot a} = \frac{1}{1} (= 1_{\mathbb{Q}})$

solve: $4 \cdot x = 1$? In \mathbb{Q} : $\frac{4}{1} \cdot x = \frac{1}{1}$ is solved by: $x = \frac{1}{4}$

Property: $(\mathbb{Q} \setminus \{0_{\mathbb{Q}}\}, \cdot)$ is an abelian group.

How to define the addition?

We want the distributive law:

$$\boxed{\frac{a}{d} + \frac{c}{d}} = \frac{a}{1} \cdot \frac{1}{d} + \frac{c}{1} \cdot \frac{1}{d} = \left(\frac{a}{1} + \frac{c}{1} \right) \cdot \frac{1}{d} = \boxed{\frac{a+c}{d}}$$

should be defined by:

$$\begin{aligned} \frac{a}{b} + \frac{c}{d} &= \frac{a}{b} \cdot \frac{d}{d} + \frac{c}{d} \cdot \frac{b}{b} = \frac{a \cdot d}{1} \cdot \frac{1}{b \cdot d} + \frac{c \cdot b}{1} \cdot \frac{1}{b \cdot d} \\ &= \left(\frac{a \cdot d}{1} + \frac{c \cdot b}{1} \right) \cdot \frac{1}{b \cdot d} = \frac{a \cdot d + c \cdot b}{b \cdot d} \end{aligned}$$

Define: $\frac{a}{b} + \frac{c}{d} := \frac{a \cdot d + c \cdot b}{b \cdot d}$ well-defined:

Proposition: The set \mathbb{Q} together with the operation $+$ and \cdot satisfies:

- (1) $(\mathbb{Q}, +)$ is an abelian group
- (2) $(\mathbb{Q} \setminus \{0_{\mathbb{Q}}\}, \cdot)$ is an abelian group
- (3) distributive law

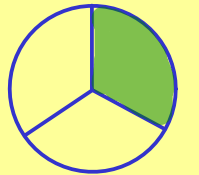
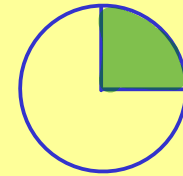
} field



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Start Learning Numbers - Part 11

$$\mathbb{Q} = \left\{ \frac{a}{b} \mid a \in \mathbb{Z}, b \in \mathbb{Z} \setminus \{0\} \right\}$$


 \leq

We need to define it



Definition of \leq for \mathbb{Z} : For $a, b \in \mathbb{Z}$, we write $a \leq b$ if $\exists k \in \mathbb{N}_0 : a + k = b$

Now: $\frac{1}{4} \leq \frac{1}{3}$ because $3 \leq 4$

Definition of \leq for \mathbb{Q} : For $b > 0$ and $d > 0$

$$\frac{a}{b} \leq \frac{c}{d} \quad \text{defined by} \quad a \cdot d \leq c \cdot b$$

Properties of \leq for \mathbb{Q} : (1) Ordering: reflexive, antisymmetric and transitive.

(2) For all $x, y, z \in \mathbb{Q}$: If $x \leq y$, then $x + z \leq y + z$

(3) For all $x, y, z \in \mathbb{Q}$: If $z \geq 0$ and $x \leq y$, then $x \cdot z \leq y \cdot z$

(4) Total order: For all $x, y \in \mathbb{Q}$, we have $x \leq y$ or $y \leq x$.

(5) Archimedean property: For all $x, \varepsilon \in \mathbb{Q}$ with $x > 0$ and $\varepsilon > 0$,

we have: $n \in \mathbb{N}_0 : n \cdot \varepsilon = \varepsilon + \varepsilon + \varepsilon + \dots + \varepsilon > x$



The Bright Side of Mathematics

The following pages cover the whole Start Learning Reals course of the Bright Side of Mathematics. Please note that the creator lives from generous supporters and would be very happy about a donation. See more here: <https://tbsom.de/support>

Have fun learning mathematics!



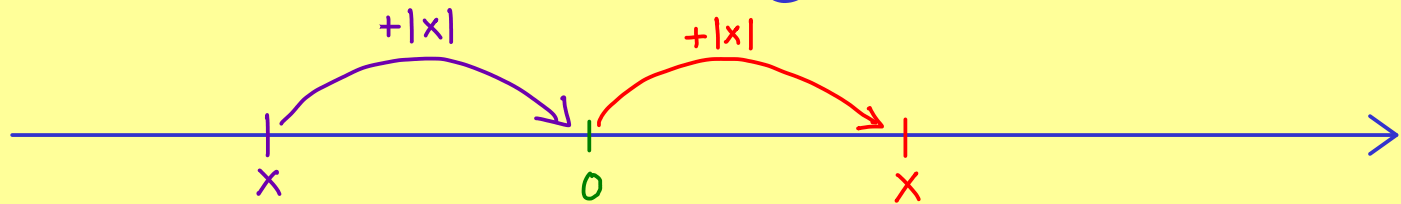
The Bright Side of Mathematics

Start Learning Reals - Part 1

} Real numbers \mathbb{R}

Starting point: \mathbb{Q} is the set of fractions \rightsquigarrow field and Archimedean order \leq
 $x > 0$, $x < 0$

Absolute value: For $x \in \mathbb{Q}$ define: $|x| := \begin{cases} x & \text{if } x \geq 0 \\ -x & \text{if } x < 0 \end{cases}$



How far away is x from 0 ? $\rightsquigarrow |x|$

Problem: There is no $x \in \mathbb{Q}$ with $x^2 = 2$

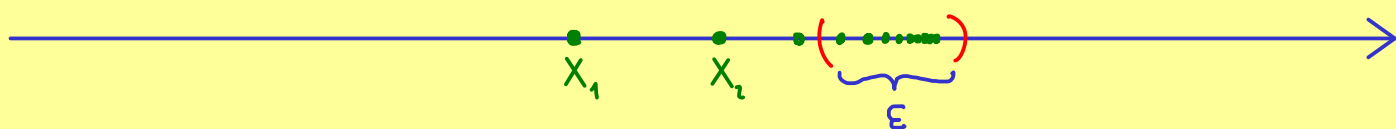
distance:
 $|x_5 - x_2|$

$x_1 = \frac{14}{10} = \frac{7}{5}$	\rightsquigarrow	$x_1^2 = \frac{49}{25} \approx 2$
$x_2 = \frac{141}{100}$	\rightsquigarrow	$x_2^2 = \frac{19881}{10000} \approx 2$
$x_3 = \frac{1414}{1000}$	\rightsquigarrow	$x_3^2 = \frac{499849}{250000} \approx 2$
$x_4 = \frac{14142}{10000}$	\rightsquigarrow	$x_4^2 = \frac{4999041}{25000000} \approx 2$
$x_5 = \frac{141421}{100000}$	\rightsquigarrow	$x_5^2 = \frac{1999899241}{10000000000} \approx 2$
\vdots	\vdots	\vdots
$x = ?$	\rightsquigarrow	$x^2 = 2$

We consider a sequence $(x_n)_{n \in \mathbb{N}}$ (infinite list; formally: a map $\mathbb{N} \rightarrow \mathbb{Q}$, $n \mapsto x_n$)
 with the property:

$$\forall \epsilon \in \mathbb{Q} \exists N \in \mathbb{N} \forall n, m \in \mathbb{N} : (\epsilon > 0 \wedge n, m \geq N \implies |x_n - x_m| < \epsilon)$$

In short: $\forall \epsilon > 0 \exists N \in \mathbb{N} \forall n, m \geq N : |x_n - x_m| < \epsilon$ (*)



Cauchy sequence: sequence $(x_n)_{n \in \mathbb{N}}$ with $x_n \in \mathbb{Q}$ and property (*)



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Start Learning Reals - Part 2

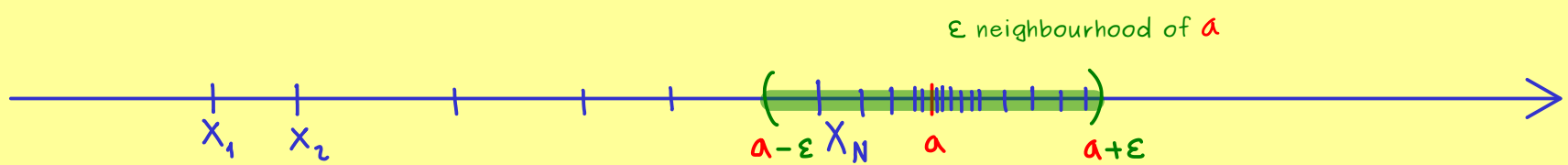
Absolute value in \mathbb{Q} : $|x \cdot y| = |x| \cdot |y|$ (multiplicative)

$|x + y| \leq |x| + |y|$ (triangle inequality)

Cauchy sequence: $(x_n)_{n \in \mathbb{N}}$ with $\forall \epsilon > 0 \exists N \in \mathbb{N} \forall n, m \geq N : |x_n - x_m| < \epsilon$

Convergent sequence: $(x_n)_{n \in \mathbb{N}}$ with $\exists a \in \mathbb{Q} \forall \epsilon > 0 \exists N \in \mathbb{N} \forall n \geq N : |x_n - a| < \epsilon$

a is called the limit of $(x_n)_{n \in \mathbb{N}}$



Example: $(\frac{1}{n})_{n \in \mathbb{N}}$ is a convergent sequence with limit $a = 0$.

Important fact: Cauchy sequence \Leftarrow Convergent sequence
not correct \mathbb{Q} but in \mathbb{R}

Proof for \Leftarrow : $|x_n - x_m| = |x_n - a + a - x_m| \leq |x_n - a| + |a - x_m|$
triangle inequality

Let $(x_n)_{n \in \mathbb{N}}$ be a convergent sequence with limit a .

Let $\epsilon > 0$. set $\epsilon' := \frac{\epsilon}{2} > 0$.

Since $(x_n)_{n \in \mathbb{N}}$ is convergent, there is $N \in \mathbb{N}$ such that:

$$\forall n \geq N : |x_n - a| < \epsilon'$$

Therefore for all $n, m \geq N$:

$$|x_n - x_m| \leq \underbrace{|x_n - a|}_{< \epsilon'} + \underbrace{|a - x_m|}_{< \epsilon'} < 2 \cdot \epsilon' = \epsilon \Rightarrow (x_n)_{n \in \mathbb{N}} \text{ Cauchy sequence}$$

Axiomatic solution: A non-empty set \mathbb{R} together with operations $+$, \cdot and ordering \leq is called the real numbers if it satisfies:

(A) $(\mathbb{R}, +, 0)$ is an abelian group

(M) $(\mathbb{R} \setminus \{0\}, \cdot, 1)$ is an abelian group ($1 \neq 0$)

(D) distributive law $x \cdot (y + z) = x \cdot y + x \cdot z$

(O) \leq is a total order, compatible with $+$ and \cdot , Archimedean property

(C) Every Cauchy sequence is a convergent sequence. $|x| := \begin{cases} x & \text{if } x \geq 0 \\ -x & \text{if } x < 0 \end{cases}$

The complete, whole, full number line \mathbb{R}

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Start Learning Reals - Part 3

complete number line \mathbb{R}

Axioms of the reals: A non-empty set \mathbb{R} together with operations $+$, \cdot and ordering \leq is called the real numbers if it satisfies:

(A) $(\mathbb{R}, +, 0)$ is an abelian group

(M) $(\mathbb{R} \setminus \{0\}, \cdot, 1)$ is an abelian group ($1 \neq 0$)

(D) distributive law $x \cdot (y + z) = x \cdot y + x \cdot z$

(O) \leq is a total order, compatible with $+$ and \cdot , Archimedean property

(C) Every Cauchy sequence is a convergent sequence. $|x| := \begin{cases} x & \text{if } x \geq 0 \\ -x & \text{if } x < 0 \end{cases}$

Important facts: There is a set with all these properties (Existence) (Construction) and it is uniquely determined by these properties. (Uniqueness) (Identification/ Isomorphism)
 ↗ see next video

Show: For all $x \in \mathbb{R}$, we have: $0 \cdot x = 0$ (*) (by only using the axioms).

$$\begin{aligned}
 \text{Proof: } 0 & \stackrel{(A)}{=} (0 \cdot x) + (-0 \cdot x) \stackrel{(A)}{=} ((0+0) \cdot x) + (-0 \cdot x) \\
 & \stackrel{(D)}{=} (0 \cdot x + 0 \cdot x) + (-0 \cdot x) \\
 & \stackrel{(A)}{=} 0 \cdot x + (0 \cdot x + (-0 \cdot x)) \stackrel{(A)}{=} 0 \cdot x + 0 \stackrel{(A)}{=} 0 \cdot x
 \end{aligned}$$

Show: For all $x \in \mathbb{R}$, we have: $(-1) \cdot x = -x$ (by only using the axioms).

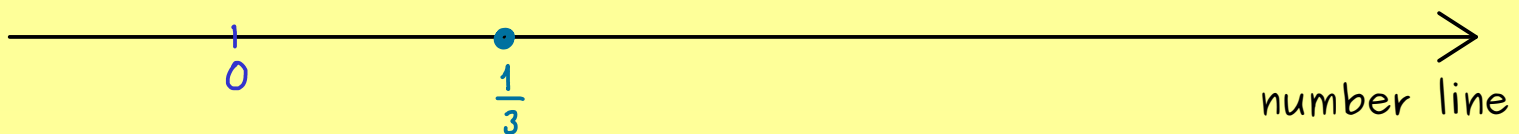
$$\begin{aligned}
 \text{Proof: } -x & \stackrel{(A)}{=} 0 + (-x) \stackrel{(*)}{=} 0 \cdot x + (-x) \stackrel{(A)}{=} ((-1)+1) \cdot x + (-x) \\
 & \stackrel{(D)}{=} (-1) \cdot x + 1 \cdot x + (-x) \stackrel{(A),(M)}{=} (-1) \cdot x + 0 \stackrel{(A)}{=} (-1) \cdot x
 \end{aligned}$$



The Bright Side of Mathematics

Start Learning Reals - Part 4

Construction: $\mathbb{Q} \rightsquigarrow \mathbb{R}$ (Make every Cauchy sequence convergent)



Sequence: $(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}, \frac{1}{3}, \frac{1}{3}, \frac{1}{3}, \dots)$ \rightsquigarrow Cauchy sequence and convergent with limit $\frac{1}{3}$

Sequence: $(\underbrace{0.3}_{\frac{3}{10}}, \underbrace{0.33}_{\frac{33}{100}}, \underbrace{0.333}_{\frac{333}{1000}}, \dots)$ \rightsquigarrow Cauchy sequence and convergent with limit $\frac{1}{3}$

$$\mathcal{C} := \left\{ (x_n)_{n \in \mathbb{N}} \mid \forall n \in \mathbb{N} : x_n \in \mathbb{Q} \text{ and } (x_n)_{n \in \mathbb{N}} \text{ is a Cauchy sequence} \right\}$$

For two elements $(a_n)_{n \in \mathbb{N}}, (b_n)_{n \in \mathbb{N}}$, define:

$$(a_n)_{n \in \mathbb{N}} \sim (b_n)_{n \in \mathbb{N}} \Leftrightarrow (a_n - b_n)_{n \in \mathbb{N}} \text{ convergent with limit } 0$$

$\Rightarrow \sim$ is an equivalence relation on \mathcal{C} (reflexive, symmetric, transitive)

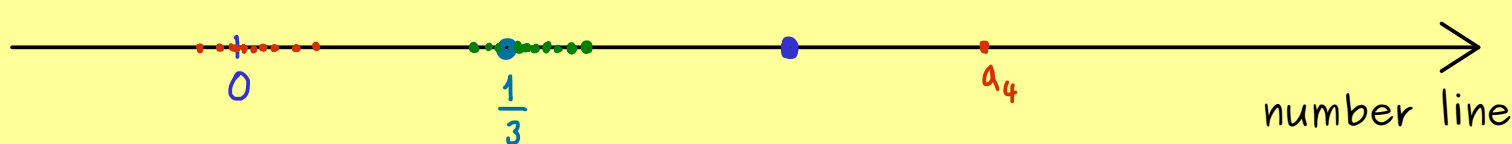
\Rightarrow equivalence class $\left[(x_n)_{n \in \mathbb{N}} \right]_{\sim} := \left\{ (a_n)_{n \in \mathbb{N}} \mid (a_n)_{n \in \mathbb{N}} \sim (x_n)_{n \in \mathbb{N}} \right\}$

Definition: $\mathbb{R} := \mathcal{C} / \sim := \left\{ \left[(x_n)_{n \in \mathbb{N}} \right]_{\sim} \mid (x_n)_{n \in \mathbb{N}} \in \mathcal{C} \right\}$

$$\left[(a_n)_{n \in \mathbb{N}} \right]_{\sim} + \left[(b_n)_{n \in \mathbb{N}} \right]_{\sim} := \left[(a_n + b_n)_{n \in \mathbb{N}} \right]_{\sim} \quad (\text{well-defined})$$

$$\left[(a_n)_{n \in \mathbb{N}} \right]_{\sim} \cdot \left[(b_n)_{n \in \mathbb{N}} \right]_{\sim} := \left[(a_n \cdot b_n)_{n \in \mathbb{N}} \right]_{\sim} \quad (\text{well-defined})$$

$$\left[(a_n)_{n \in \mathbb{N}} \right]_{\sim} < \left[(b_n)_{n \in \mathbb{N}} \right]_{\sim} \Leftrightarrow \exists \delta > 0 \exists N \in \mathbb{N} \forall n \geq N : \delta < b_n - a_n$$



Properties: (A) $(\mathbb{R}, +, 0)$ is an abelian group

(M) $(\mathbb{R} \setminus \{0\}, \cdot, 1)$ is an abelian group ($1 \neq 0$)

(D) distributive law $x \cdot (y + z) = x \cdot y + x \cdot z$

(O) \leq is a total order, compatible with $+$ and \cdot , Archimedean property

(C) Every Cauchy sequence is a convergent sequence. $|x| := \begin{cases} x & \text{if } x \geq 0 \\ -x & \text{if } x < 0 \end{cases}$

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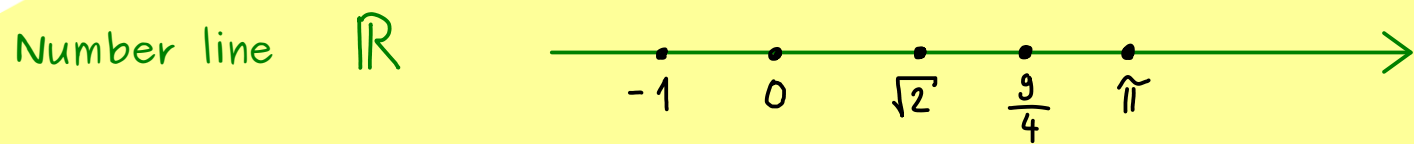
The following pages cover the whole Start Learning Complex Numbers course of the Bright Side of Mathematics. Please note that the creator lives from generous supporters and would be very happy about a donation. See more here: <https://tbsom.de/support>

Have fun learning mathematics!



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Start Learning Complex Numbers - Part 1



- field $+$ \cdot
- ordering \leq (Archimedean, compatible with $+$ and \cdot , ...)
- complete

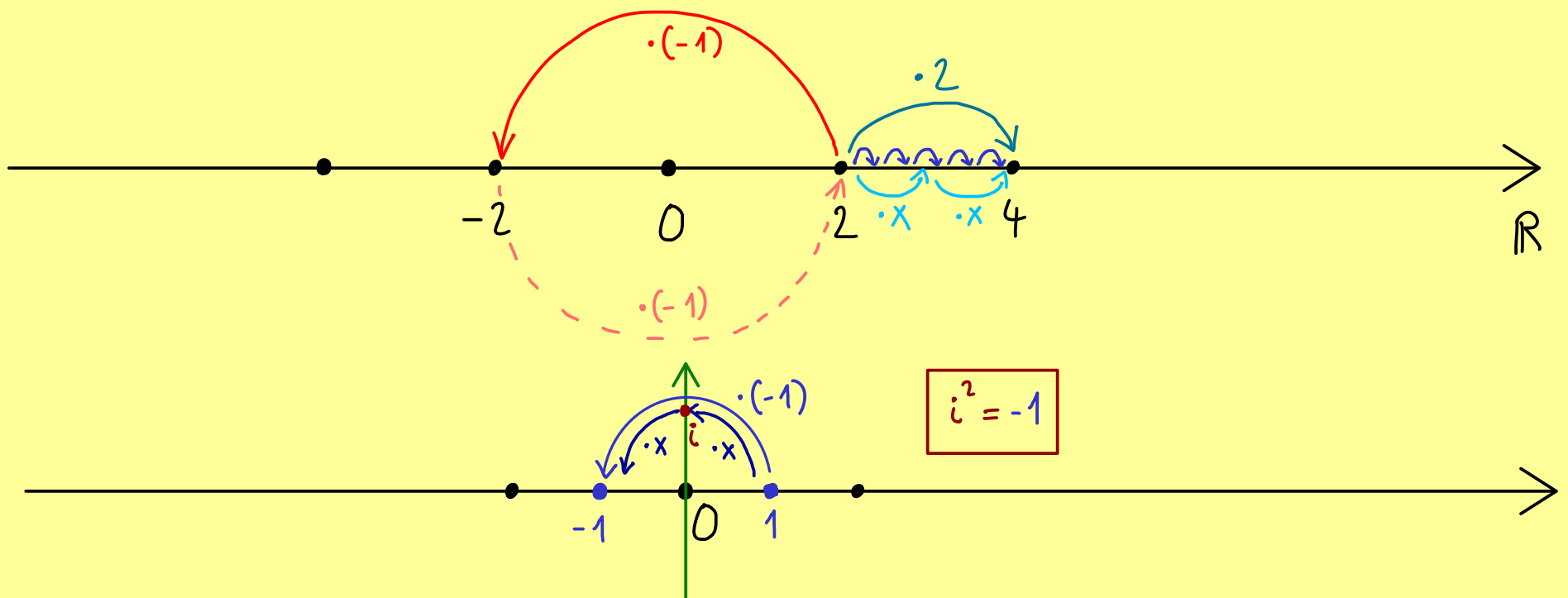
One can solve a lot of equations:

$$x + 5 = 1, \quad x + x = -1$$

$$x \cdot 5 = 1$$

$$x^2 = 2$$

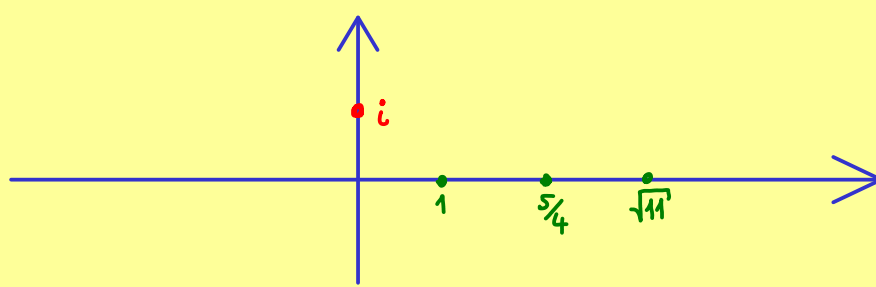
We cannot solve: $x^2 = -1$ (because $x^2 \geq 0$ for all $x \in \mathbb{R}$ and $-1 < 0$)





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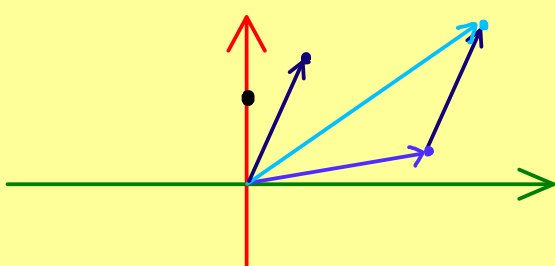
Start Learning Complex Numbers - Part 2



$$\mathbb{R} \times \mathbb{R} = \mathbb{R}^2$$

+ addition
• multiplication

Addition: For $\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}, \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} \in \mathbb{R} \times \mathbb{R}$, we set: $\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} + \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} := \begin{pmatrix} x_1 + y_1 \\ x_2 + y_2 \end{pmatrix}$



Multiplication: For $\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}, \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} \in \mathbb{R} \times \mathbb{R}$, we set: $\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} \cdot \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} := \begin{pmatrix} x_1 y_1 - x_2 y_2 \\ x_2 y_1 + x_1 y_2 \end{pmatrix}$

Why?

Short notation: $\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} := x_1 + i \cdot x_2$, $\begin{pmatrix} 0 \\ 1 \end{pmatrix} = 0 + i \cdot 1 := i$

$$\begin{aligned} \text{Calculation: } (x_1 + i \cdot x_2) \cdot (y_1 + i \cdot y_2) &= x_1 y_1 + i \cdot x_2 y_1 + i \cdot x_1 y_2 + i^2 x_2 y_2 \\ &= (x_1 y_1 - x_2 y_2) + i \cdot (x_2 y_1 + x_1 y_2) \end{aligned}$$

Check: $i^2 = i \cdot i = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \cdot \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 - 1 \\ 0 + 0 \end{pmatrix} = \begin{pmatrix} -1 \\ 0 \end{pmatrix} = -1 + i \cdot 0 = -1$

Properties: • We write $\mathbb{C} := \mathbb{R}^2$ when we have + and • from above.

- field {
- $(\mathbb{C}, +, 0)$ is an abelian group (commutative, associative, neutral element, inverses)
 $\stackrel{=}{=} 0 + i \cdot 0$
 - $(\mathbb{C} \setminus \{0\}, \cdot, 1)$ is an abelian group (commutative, associative, neutral element, inverses)
 $\stackrel{=}{=} 1 + i \cdot 0$
 - distributive law
 - no nice ordering \leq like for \mathbb{R}

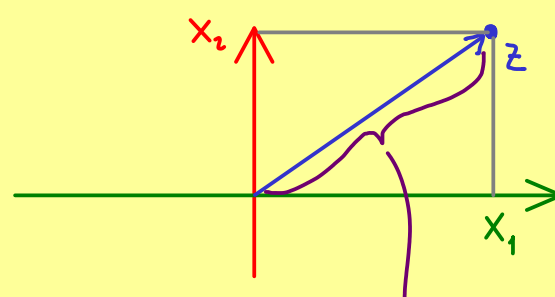


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Start Learning Complex Numbers - Part 3

$$z = x_1 + i \cdot x_2 \in \mathbb{C}$$

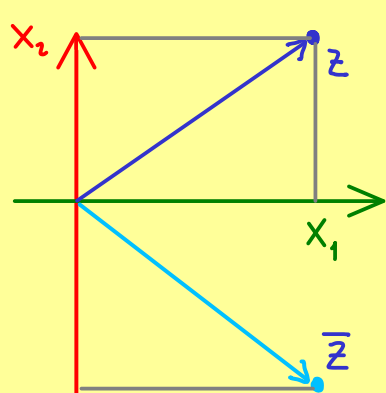
real part of z $\text{Re}(z)$
 imaginary part of z $\text{Im}(z)$



length, absolute value, modulus

$$|z| := \sqrt{\text{Re}(z)^2 + \text{Im}(z)^2} \in \mathbb{R}$$

Reflection:
complex conjugate



$$z = x_1 + i \cdot x_2$$

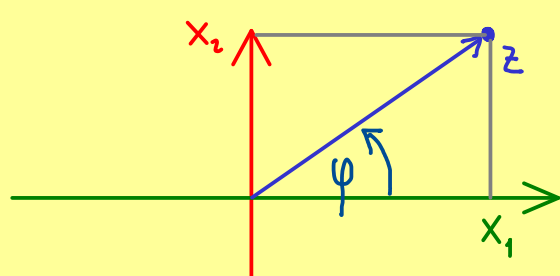
$$\bar{z} = x_1 + i \cdot (-x_2) = x_1 - i \cdot x_2$$

Calculate: $z \cdot \bar{z} = (x_1 + i \cdot x_2) \cdot (x_1 - i \cdot x_2)$

$$= x_1^2 + x_1 \cdot (-i \cdot x_2) + i \cdot x_2 \cdot x_1 - i^2 \cdot x_2^2$$

$$= x_1^2 + x_2^2 = |z|^2$$

Polar coordinates:



length: $|z|$

angle: $\varphi \in [0, 2\pi)$

argument of z :

$$\varphi = \begin{cases} \arctan\left(\frac{x_2}{x_1}\right) & , x_1 > 0, x_2 \geq 0 \\ \frac{\pi}{2} & , x_1 = 0, x_2 > 0 \\ \arctan\left(\frac{x_2}{x_1}\right) + \pi & , x_1 < 0 \\ \frac{3\pi}{2} & , x_1 = 0, x_2 < 0 \\ \arctan\left(\frac{x_2}{x_1}\right) + 2\pi & , x_1 > 0, x_2 < 0 \end{cases}$$

$$z = x_1 + i \cdot x_2 = |z| \cdot (\cos(\varphi) + i \cdot \sin(\varphi))$$

Example: $z = 3 + i \cdot 3$, $\bar{z} = 3 - i \cdot 3$, $z \cdot \bar{z} = 9 + 9 = 18$

$$\Rightarrow |z| = \sqrt{18} = 3 \cdot \sqrt{2} \quad , \quad \varphi = \arctan\left(\frac{3}{3}\right) = \frac{\pi}{4}$$

$$\Rightarrow z = 3 \cdot \sqrt{2} \cdot (\cos(\frac{\pi}{4}) + i \cdot \sin(\frac{\pi}{4})) \stackrel{\text{later}}{=} 3 \cdot \sqrt{2} \cdot e^{i \frac{\pi}{4}}$$